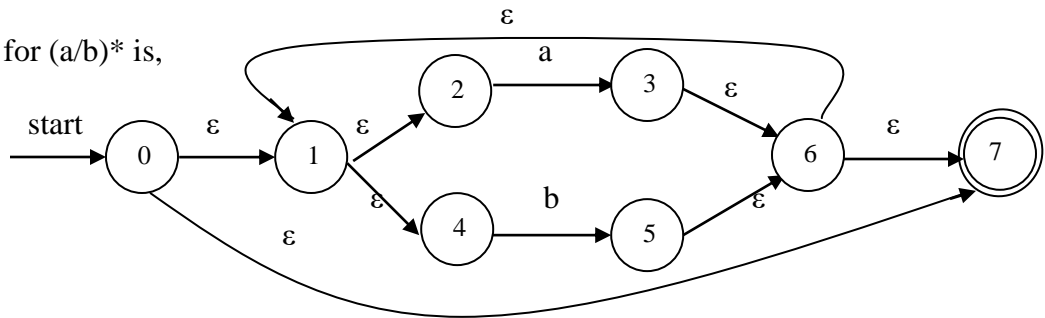


## CONVERSION OF NFA INTO DFA

1. Convert the NFA  $(a/b)^*$  into DFA?

Solution:

The NFA for  $(a/b)^*$  is,



$\epsilon$  closure  $\{0\} = \{0,1,2,4,7\}$  ----- A

Transition of input symbol a on A =  $\{3\}$

Transition of input symbol b on A =  $\{5\}$

$\epsilon$  closure  $\{3\} = \{3,6,1,2,4,7\}$  ----- B

Transition of input symbol a on B =  $\{3\}$

Transition of input symbol b on B =  $\{5\}$

$\epsilon$  closure  $\{5\} = \{5,6,1,2,4,7\}$  ----- C

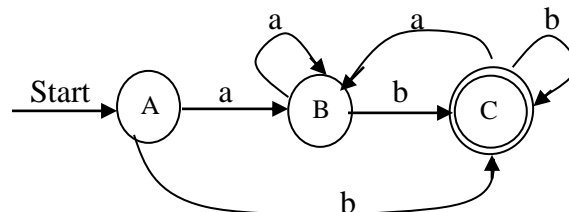
Transition of input symbol a on C =  $\{3\}$

Transition of input symbol b on C =  $\{5\}$

Since A is the start state and state C is the only accepting state then, the transition table is,

| State | Input symbol |   |
|-------|--------------|---|
|       | a            | b |
| A     | B            | C |
| B     | B            | C |
| C     | B            | C |

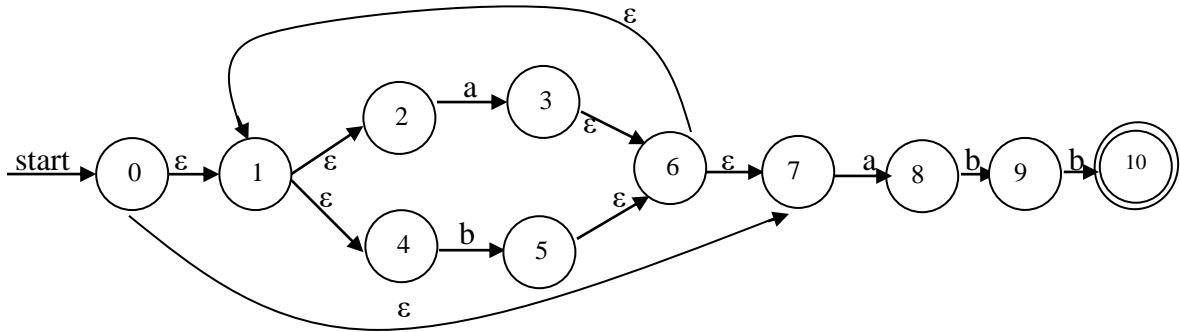
The DFA is,



2. Convert the NFA  $(a/b)^*abb$  into DFA?

Solution:

The NFA for  $(a/b)^*abb$  is,



$\epsilon$  closure  $\{0\} = \{0,1,2,4,7\}$  ----- A

Transition of input symbol a on A =  $\{3,8\}$

Transition of input symbol b on A =  $\{5\}$

$\epsilon$  closure  $\{3,8\} = \{3,6,7,1,2,4,8\}$  ----- B

Transition of input symbol a on B =  $\{8,3\}$

Transition of input symbol b on B =  $\{5,9\}$

$\epsilon$  closure  $\{5\} = \{5,6,7,1,2,4\}$  ----- C

Transition of input symbol a on C =  $\{8,3\}$

Transition of input symbol b on C =  $\{5\}$

$\epsilon$  closure  $\{5,9\} = \{5,6,7,1,2,4,9\}$  ----- D

Transition of input symbol a on D =  $\{8,3\}$

Transition of input symbol b on D =  $\{5,10\}$

$\epsilon$  closure  $\{5,10\} = \{5,6,7,1,2,4,10\}$  ----- E

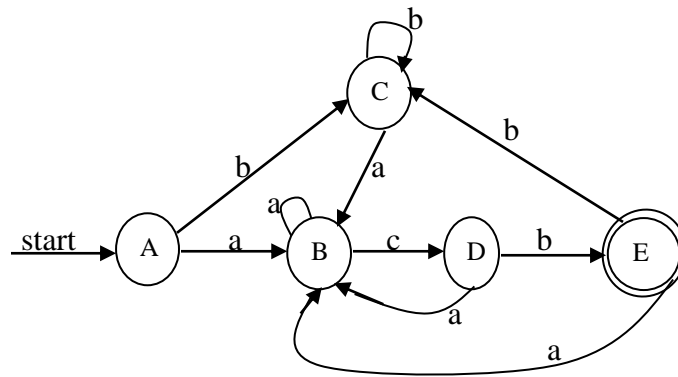
Transition of input symbol a on E =  $\{8,3\}$

Transition of input symbol b on E =  $\{5\}$

Since A is the start state and state E is the only accepting state then, the **transition table** is,

| State | Input symbol |   |
|-------|--------------|---|
|       | a            | b |
| A     | B            | C |
| B     | B            | D |
| C     | B            | C |
| D     | B            | E |
| E     | B            | C |



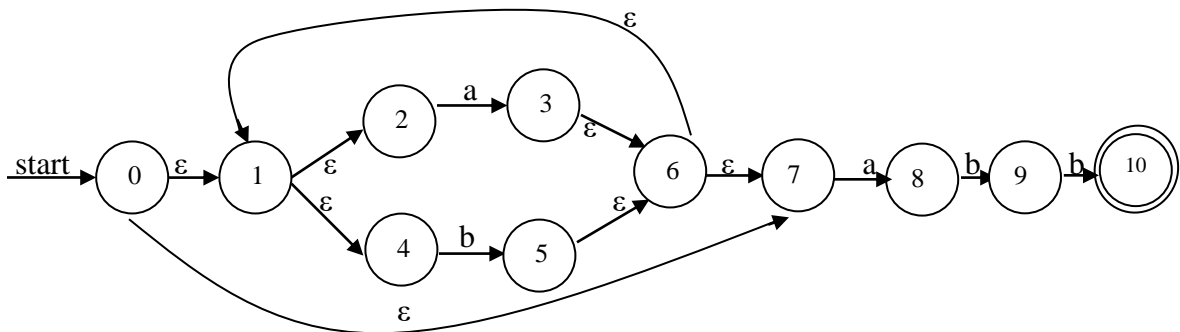


**MINIMIZATION OF STATES**

Problem 1: Construct a minimum state DFA for a regular expression  $(a/b)^*abb$

Solution:-

1. The NFA of  $(a/b)^*abb$  is



2. Construct a DFA:

$\epsilon$  closure  $\{0\} = \{0,1,2,4,7\}$  ----- A

Transition of input symbol a on A =  $\{3,8\}$

Transition of input symbol b on A =  $\{5\}$

$\epsilon$  closure  $\{3,8\} = \{3,6,7,1,2,4,8\}$  ----- B

Transition of input symbol a on B =  $\{8,3\}$

Transition of input symbol b on B =  $\{5,9\}$

$\epsilon$  closure  $\{5\} = \{5,6,7,1,2,4\}$  ----- C

Transition of input symbol a on C =  $\{8,3\}$

Transition of input symbol b on C = { 5 }

$\epsilon$  closure {5,9} = { 5,6,7,1,2,4,9} ----- D

Transition of input symbol a on D = { 8,3 }

Transition of input symbol b on D = { 5,10 }

$\epsilon$  closure {5,10} = { 5,6,7,1,2,4,10} ----- E

Transition of input symbol a on E = { 8,3 }

Transition of input symbol b on E = { 5 }

Since A is the start state and state E is the only accepting state then, the **transition table** is,

| State | Input symbol |   |
|-------|--------------|---|
|       | a            | b |
| A     | B            | C |
| B     | B            | D |
| C     | B            | C |
| D     | B            | E |
| E     | B            | C |

### 3. Minimizing the DFA

Let  $\Pi = ABCDE$

The initial partition  $\Pi$  consists of two groups.

$\Pi_1 = ABCD$  ( that is the non – accepting states)

$\Pi_2 = E$  ( that is the accepting state)

So, (ABCD) (E)

#### AB

A  $\xrightarrow{a}$  B      B  $\xrightarrow{a}$  B

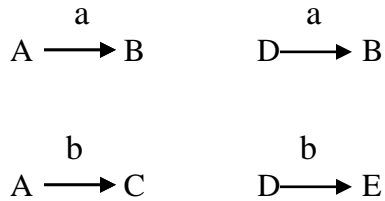
A  $\xrightarrow{b}$  C      B  $\xrightarrow{b}$  D

#### AC

A  $\xrightarrow{a}$  B      C  $\xrightarrow{a}$  B

A  $\xrightarrow{b}$  C      C  $\xrightarrow{b}$  C

AD

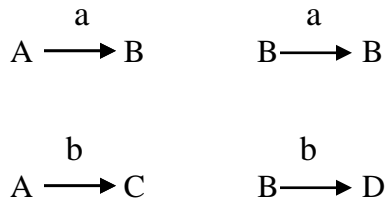


On input "a" each of these states has a transition to B, so they could all remain in one group as far as input a is concerned.

On input "b" A,B,C go to members of the group  $\Pi_1$  (ABC) while D goes to  $\Pi_2$  (E). Thus  $\Pi_1$  group is split into two new groups.

$\Pi_1 = ABC$      $\Pi_2 = D$ ,  $\Pi_3 = E$   
So, (ABC) (D) (E)

AB

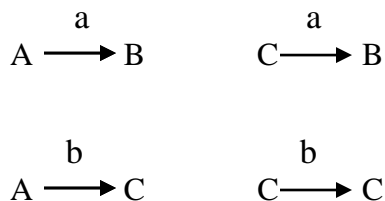


Here B goes to  $\Pi_2$ . Thus  $\Pi_1$  group is again split into two new groups. The new groups are,

$\Pi_1 = AC$      $\Pi_2 = B$ ,  $\Pi_3 = D$ ,  $\Pi_4 = E$   
So, (AC) (B) (D) (E)

Here we cannot split any of the groups consisting of the single state. The only possibility is try to split only (AC)

For AC



But A and C go the same state B on input a, and they go to the same state C on input b.

Hence after this,

(AC) (B) (D) (E)

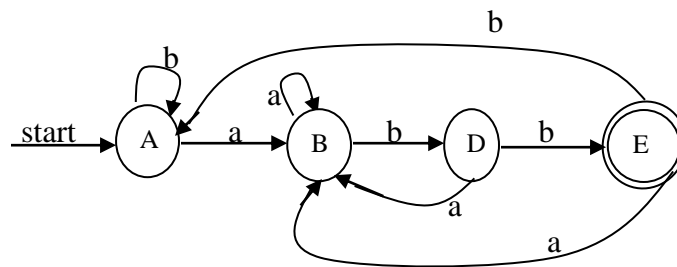
Here we choose A as the representative for the group AC.

Thus A is the start state and state E is the only accepting state.

So the *minimized transition table* is,

| State | Input symbol |   |
|-------|--------------|---|
|       | a            | b |
| A     | B            | A |
| B     | B            | D |
| D     | B            | E |
| E     | B            | A |

Thus the minimized DFA is,



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