# Lecture A3: <br> Finite Fields of the Form GF(2 $\mathbf{2}^{\text {n }}$ ) 

$4^{\text {th }}$ Year Course- CCSIT, UoA

## Lecture goals

$\square$ To review finite fields of the form GF (2n)
$\square$ To show how arithmetic operations can be carried out by directly operating on the bit patterns for the elements of GF ( $2^{n}$ )

## Consider Again the Polynomials Over GF (2)

- Here are some examples:


```
x
1
x
x}1000
```

- We could also shown polynomials with negative coefficients, but recall that in $G F(2),-1$ is the same as +1 .
- Obviously, the number of such polynomials is infinite.
- The polynomials can be subject to the algebraic operations of addition and multiplication in which the coefficients are added and multiplied according to the rules that apply to GF (2).
- As stated in the previous lecture, the set of such polynomials forms a ring, called the polynomial ring.


## Modular Polynomial Arithmetic (1)

$>$ Let's now add one more twist to the algebraic operations we carry out on all the polynomials over GF (2):
$>$ We will first choose a particular irreducible polynomial, as for example $\quad x^{3}+x+1$
(By the way there exist only two irreducible polynomials of degree 3 over GF (2). The other is $x^{3}+x^{2}+1$.)
$>$ We will now consider all polynomials defined over GF (2) modulo the irreducible polynomial $x^{3}+x+1$.
$>$ In particular, when an algebraic operation (we are obviously talking about polynomial multiplication) results in a polynomial whose degree equals or exceeds that of the irreducible polynomial, we will take for our result the remainder modulo the irreducible polynomial.

## Modular Polynomial Arithmetic

$>$ For example,

$$
\begin{aligned}
& \left(x^{2}+x+1\right) \times\left(x^{2}+1\right) \bmod \left(x^{3}+x+1\right) \\
& \quad=\left(x^{4}+x^{3}+x^{2}\right)+\left(x^{2}+x+1\right) \bmod \left(x^{3}+x+1\right) \\
& =\left(x^{4}+x^{3}+x+1\right) \bmod \left(x^{3}+x+1\right) \\
& =-x^{2}-x \\
& =x^{2}+x
\end{aligned}
$$

$>$ Recall that $1+1=0$ in GF (2). This is what we used in getting to the second expression on the right hand side.
$>$ For the division by the modulus in the above example, we used the result

$$
\frac{\left(x^{4}+x^{3}+x+1\right)}{\left(x^{3}+x+1\right)}=x+1+\frac{-x^{2}-x}{x^{3}+x+1}
$$

$>$ Obviously, for the division on the left hand side, our first quotient term is $x$. Multiplying the divisor by $x$ yields $x^{4}+x^{2}+x$ that when subtracted from the dividend gives us $x^{3}-x^{2}+1$. This dictates that the next term of the quotient be 1 , and so on.

## How Large is the Set of Polynomials When Multiplications are Carried Out Modulo $x^{3}+x+1$

$\square$ With multiplications modulo $x^{3}+x+1$, we have only the following eight polynomials in the set of polynomials over GF (2):


We will refer to this set as GF $\left(2^{3}\right)$ where the power of 2 is the degree of the modulus polynomial.
Our conceptualization of GF $\left(2^{3}\right)$ is analogous to our conceptualization of the set $Z_{8}$. The eight elements of $Z_{8}$ are to be thought of as integers modulo 8. So, basically, $Z_{8}$ maps all integers to the eight in the set $Z_{8}$. Similarly, $G F\left(2^{3}\right)$ maps all of the polynomials over $G F(2)$ to the eight polynomials shown above.
But note the crucial difference between $\operatorname{GF}\left(2^{3}\right)$ and $Z_{8}: G F\left(2^{3}\right)$ is a field, whereas $Z_{8}$ is NOT.

* We do know that $G F\left(2^{3}\right)$ is an abelian group because of the operation of polynomial addition satisfies all of the requirements on a group operator and because polynomial addition is commutative.
* GF $\left(2^{3}\right)$ is also a commutative ring because polynomial multiplication distributes over polynomial addition (and because polynomial multiplication meets all the other stipulations on the ring operator: closedness, associativity, commutativity).
GF $\left(2^{3}\right)$ is an integral domain because of the fact that the set contains the multiplicative identity element 1 and because if for $a \in G F\left(2^{3}\right)$ and $b \in G F\left(2^{3}\right)$ we have

$$
a \times b=0 \bmod \left(x^{3}+x+1\right)
$$

then either $a=0$ or $b=0$.

* $G F\left(2^{3}\right)$ is a finite field because it is a finite set and because it contains a unique multiplicative inverse for every non-zero element.
*GF $\left(2^{3}\right)$ contains a unique multiplicative inverse for every non- zero element for the same reason that $Z_{7}$ contains a unique multiplicative inverse for every non-zero integer in the set. (For a counterexample, recall that $Z_{8}$ does not possess multiplicative inverses for 2,4 , and 6 .)
* In other words, for every non-zero element $a \in G F\left(2^{3}\right)$ there is always a unique element $b \in G F\left(2^{3}\right)$ such that $a \times b=1$.
This follows from the fact if you multiply a non-zero element $a$ with each of the eight elements of $G F\left(2^{3}\right)$, the result will the eight distinct elements of $G F\left(2^{3}\right)$.


## How Do We Know That $G F\left(2^{3}\right)$ is a Finite Field?

Obviously, the results of such multiplications must equal 1 for exactly one of the non- zero element of $\operatorname{GF}\left(2^{3}\right)$. So if $a \times$ $b=1$, then $b$ must be the multiplicative inverse for $a$.

* The same thing happens in $Z_{7}$. If you multiply a non-zero element $a$ of this set with each of the seven elements of $Z_{7}$, you will get seven distinct answers. The answer must therefore equal 1 for at least one such multiplication. When the answer is 1 , you have your multiplicative inverse for $a$.
For a counterexample, this is not what happens in $Z_{8}$. When you multiply 2 with every element of $Z_{8}$, you do not get eight distinct answers. (Multiplying 2 with every element of $Z_{8}$ yields $\{0,2,4,6,0,2,4,6\}$ that has only four distinct elements).

The upshot is that $G F\left(2^{3}\right)$ is a finite field.

## $G F\left(2^{n}\right)$ is a Finite Field for Every $n$

$>$ None of the arguments on the previous three pages is limited by the value 3 for the power of 2 . That means that $G F\left(2^{n}\right)$ is a finite field for every $n$.
$>$ To find all the polynomials in GF $\left(2^{n}\right)$, we obviously need an irreducible polynomial of degree $n$.
$>$ AES arithmetic is based on GF $\left(2^{8}\right)$. It uses the following irreducible polynomial

$$
x^{8}+x^{4}+x^{3}+x+1
$$

$>$ The finite field $G F\left(2^{8}\right)$ used by AES obviously contains 256 distinct polynomials over GF (2).
$>$ In general, $G F\left(p^{n}\right)$ is a finite field for any prime $p$. The elements of $G F\left(p^{n}\right)$ are polynomials over $G F(p)$ (which is the same as the set of residues $Z_{p}$ ).

## Representing the Individual Polynomials in $G F\left(2^{n}\right)$ by Binary Code Words

$\square$ Recall the eight polynomials in $G F\left(2^{3}\right)$ when the modulus polynomial is $x^{3}+x+1$ (See the next page).
$\square$ We now claim that there is nothing sacred about the variable $x$ in such polynomials.
$\square$ We can think of $x^{i}$ as being merely a place-holder for a bit.
$\square$ That is, we can think of the polynomials as bit strings corresponding to the coefficients that can only be 0 or 1 , each power of $x$ representing a specific position in a bit string.
$\square$ So the $2^{3}$ polynomials of $G F\left(2^{3}\right)$ can therefore be represented by the bit strings shown in the next page.

## Representing the Individual Polynomials in $G F\left(2^{n}\right)$ by Binary Code Words

| 0 | $\Rightarrow$ | 000 |
| ---: | :--- | :--- |
| 1 | $\Rightarrow$ | 001 |
| $x$ | $\Rightarrow$ | 010 |
| $x+1$ | $\Rightarrow$ | 011 |
| $x^{2}$ | $\Rightarrow$ | 100 |
| $x^{2}+1$ | $\Rightarrow$ | 101 |
| $x^{2}+x$ | $\Rightarrow$ | 110 |
| $x^{2}+x+1$ |  | 111 |

$\square$ If we wish, we can give a decimal representation to each of the above bit patterns. The decimal values between 0 and 7, both limits inclusive, would have to obey the addition and multiplication rules corresponding to the underlying finite field.
$\square$ Exactly the same approach can be used to come up with $2^{n}$ bit patterns, each pattern consisting of $n$ bits, for a set of integers that would constitute a finite field, provided we have available to us an irreducible polynomial of degree $n$.

