## Chapter 2

## Solutions of Equations of One Variable

### 2.1 Introduction

In this chapter we consider one of the most basic problems of numerical approximation, the root-finding problem. This process involves finding a root, or solution, of an equation of the form $f(x)=0$. A root of this equation is also called a zero of the function $f$. This is one of the oldest known approximation problems, yet research continues in this area at the present time.

The problem of finding an approximation to the root of an equation can be traced at least as far back as 1700 B.c. A cuneiform table in the Yale Babylonian Collection dating from that period gives a sexagesimal (base-60) number equivalent to 1.414222 as an approximation to $\sqrt{2}$, a result that is accurate to within $10^{-5}$. This approximation can be found by applying a technique given in Section 2.4.

### 2.2 The Bisection Method

The first and most elementary technique we consider is the Bisection, or BinarySearch, method. The Bisection method is used to determine, to any specified accuracy that your computer will permit, a solution to $f(x)=0$ on an interval $[a, b]$, provided that $f$ is continuous on the interval and that $f(a)$ and $f(b)$ are of opposite sign. Although the method will work for the case when more than one root is contained in the interval $[a, b]$, we assume for simplicity of our discussion that the root in this interval is unique.

To begin the Bisection method, set $a_{1}=a$ and $b_{1}=b$, as shown in Figure 2.1, and let $p_{1}$ be the midpoint of the interval $[a, b]$ :

$$
p_{1}=a_{1}+\frac{b_{1}-a_{1}}{2}
$$

If $f\left(p_{1}\right)=0$, then the root $p$ is given by $p=p_{1}$; if $f\left(p_{1}\right) \neq 0$, then $f\left(p_{1}\right)$ has the same sign as either $f\left(a_{1}\right)$ or $f\left(b_{1}\right)$.

## Figure 2.1



If $f\left(p_{1}\right)$ and $f\left(a_{1}\right)$ have the same sign, then $p$ is in the interval $\left(p_{1}, b_{1}\right)$, and we set

$$
a_{2}=p_{1} \quad \text { and } \quad b_{2}=b_{1}
$$

If, on the other hand, $f\left(p_{1}\right)$ and $f\left(a_{1}\right)$ have opposite signs, then $p$ is in the interval $\left(a_{1}, p_{1}\right)$, and we set

$$
a_{2}=a_{1} \quad \text { and } \quad b_{2}=p_{1}
$$

We reapply the process to the interval $\left[a_{2}, b_{2}\right]$, and continue forming $\left[a_{3}, b_{3}\right]$, $\left[a_{4}, b_{4}\right], \ldots$ Each new interval will contain $p$ and have length one half of the length of the preceding interval.
[Bisection Method] An interval $\left[a_{n+1}, b_{n+1}\right]$ containing an approximation to a root of $f(x)=0$ is constructed from an interval $\left[a_{n}, b_{n}\right]$ containing the root by first letting

$$
p_{n}=a_{n}+\frac{b_{n}-a_{n}}{2}
$$

Then set

$$
a_{n+1}=a_{n} \quad \text { and } \quad b_{n+1}=p_{n} \quad \text { if } \quad f\left(a_{n}\right) f\left(p_{n}\right)<0
$$

and

$$
a_{n+1}=p_{n} \quad \text { and } \quad b_{n+1}=b_{n} \quad \text { otherwise }
$$

There are three stopping criteria commonly incorporated in the Bisection method. First, the method stops if one of the midpoints happens to coincide with the root. It also stops when the length of the search interval is less than some prescribed tolerance we call $T O L$. The procedure also stops if the number of iterations exceeds a preset bound $N_{0}$.

To start the Bisection method, an interval $[a, b]$ must be found with $f(a) \cdot f(b)<$ 0 . At each step, the length of the interval known to contain a zero of $f$ is reduced by a factor of 2 . Since the midpoint $p_{1}$ must be within $(b-a) / 2$ of the root $p$, and each succeeding iteration divides the interval under consideration by 2 , we have

$$
\left|p_{n}-p\right| \leq \frac{b-a}{2^{n}}
$$

As a consequence, it is easy to determine a bound for the number of iterations needed to ensure a given tolerance. If the root needs to be determined within the tolerance $T O L$, we need to determine the number of iterations, $n$, so that

$$
\frac{b-a}{2^{n}}<T O L
$$

Solving for $n$ in this inequality gives

$$
\frac{b-a}{T O L}<2^{n}, \quad \text { which implies that } \quad \log _{2}\left(\frac{b-a}{\mathrm{TOL}}\right)<n .
$$

Since the number of required iterations to guarantee a given accuracy depends on the length of the initial interval $[a, b]$, we want to choose this interval as small as possible. For example, if $f(x)=2 x^{3}-x^{2}+x-1$, we have both

$$
f(-4) \cdot f(4)<0 \quad \text { and } \quad f(0) \cdot f(1)<0
$$

so the Bisection method could be used on either $[-4,4]$ or $[0,1]$. Starting the $\mathrm{Bi}-$ section method on $[0,1]$ instead of $[-4,4]$ reduces by 3 the number of iterations required to achieve a specified accuracy.

EXAMPLE 1 The equation $f(x)=x^{3}+4 x^{2}-10=0$ has a root in $[1,2]$ since $f(1)=-5$ and $f(2)=14$. It is easily seen from a sketch of the graph of $f$ in Figure 2.2 that there is only one root in $[1,2]$.

Figure 2.2


To use Maple to approximate the root, we define the function $f$ by the command
$>f:=x->x^{\wedge} 3+4 * x^{\wedge} 2-10 ;$

The values of $a_{1}$ and $b_{1}$ are given by
>a1:=1; b1:=2;

We next compute $f\left(a_{1}\right)=-5$ and $f\left(b_{1}\right)=14$ by

```
>fa1:=f(a1); fb1:=f(b1);
```

and the midpoint $p_{1}=1.5$ and $f\left(p_{1}\right)=2.375$ by

```
>p1:=a1+0.5*(b1-a1);
>pf1:=f(p1);
```

Since $f\left(a_{1}\right)$ and $f\left(p_{1}\right)$ have opposite signs, we reject $b_{1}$ and let $a_{2}=a_{1}$ and $b_{2}=p_{1}$. This process is continued to find $p_{2}, p_{3}$, and so on.

As discussed in the Preface, each of the methods we consider in the book has an accompanying set of programs contained on the CD that is in the back of the book. The programs are given for the programming languages C, FORTRAN, and Pascal, and also in Maple V, Mathematica, and MATLAB. The program BISECT21, provided with the inputs $a=1, b=2, T O L=0.0005$, and $N_{0}=20$, gives the values in Table 2.1. The actual root $p$, to 10 decimal places, is $p=1.3652300134$, and $\left|p-p_{11}\right|<0.0005$. Since the expected number of iterations is $\log _{2}((2-1) / 0.0005) \approx$ 10.96, the bound $N_{0}$ was certainly sufficient.

Table 2.1

| $n$ | $a_{n}$ | $b_{n}$ | $p_{n}$ | $f\left(p_{n}\right)$ |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 1.0000000000 | 2.0000000000 | 1.5000000000 | 2.3750000000 |
| 2 | 1.0000000000 | 1.5000000000 | 1.2500000000 | -1.7968750000 |
| 3 | 1.2500000000 | 1.5000000000 | 1.3750000000 | 0.1621093750 |
| 4 | 1.2500000000 | 1.3750000000 | 1.3125000000 | -0.8483886719 |
| 5 | 1.3125000000 | 1.3750000000 | 1.3437500000 | -0.3509826660 |
| 6 | 1.3437500000 | 1.3750000000 | 1.3593750000 | -0.0964088440 |
| 7 | 1.3593750000 | 1.3750000000 | 1.3671875000 | 0.0323557854 |
| 8 | 1.3593750000 | 1.3671875000 | 1.3632812500 | -0.0321499705 |
| 9 | 1.3632812500 | 1.3671875000 | 1.3652343750 | 0.0000720248 |
| 10 | 1.3632812500 | 1.3652343750 | 1.3642578125 | -0.0160466908 |
| 11 | 1.3642578125 | 1.3652343750 | 1.3647460938 | -0.0079892628 |

The Bisection method, though conceptually clear, has serious drawbacks. It is slow to converge relative to the other techniques we will discuss, and a good intermediate approximation may be inadvertently discarded. This happened, for example, with $p_{9}$ in Example 1. However, the method has the important property that it always converges to a solution and it is easy to determine a bound for the number of iterations needed to ensure a given accuracy. For these reasons, the Bisection method is frequently used as a dependable starting procedure for the more efficient methods presented later in this chapter.

The bound for the number of iterations for the Bisection method assumes that the calculations are performed using infinite-digit arithmetic. When implementing the method on a computer, consideration must be given to the effects of round-off error. For example, the computation of the midpoint of the interval $\left[a_{n}, b_{n}\right]$ should be found from the equation

$$
p_{n}=a_{n}+\frac{b_{n}-a_{n}}{2}
$$

instead of from the algebraically equivalent equation

$$
p_{n}=\frac{a_{n}+b_{n}}{2}
$$

The first equation adds a small correction, $\left(b_{n}-a_{n}\right) / 2$, to the known value $a_{n}$. When $b_{n}-a_{n}$ is near the maximum precision of the machine, this correction might be in error, but the error would not significantly affect the computed value of $p_{n}$. However, in the second equation, if $b_{n}-a_{n}$ is near the maximum precision of the machine, it is possible for $p_{n}$ to return a midpoint that is not even in the interval [ $a_{n}, b_{n}$ ].

A number of tests can be used to see if a root has been found. We would normally use a test of the form

$$
\left|f\left(p_{n}\right)\right|<\varepsilon
$$

where $\varepsilon>0$ would be a small number related in some way to the tolerance. However, it is also possible for the value $f\left(p_{n}\right)$ to be small when $p_{n}$ is not near the root $p$.

As a final remark, to determine which subinterval of $\left[a_{n}, b_{n}\right]$ contains a root of $f$, it is better to make use of signum function, which is defined as

$$
\operatorname{sgn}(x)= \begin{cases}-1, & \text { if } x<0 \\ 0, & \text { if } x=0 \\ 1, & \text { if } x>0\end{cases}
$$

The test

$$
\operatorname{sgn}\left(f\left(a_{n}\right)\right) \operatorname{sgn}\left(f\left(p_{n}\right)\right)<0 \quad \text { instead of } \quad f\left(a_{n}\right) f\left(p_{n}\right)<0
$$

gives the same result but avoids the possibility of overflow or underflow in the multiplication of $f\left(a_{n}\right)$ and $f\left(p_{n}\right)$.

## EXERCISE SET 2.2

1. Use the Bisection method to find $p_{3}$ for $f(x)=\sqrt{x}-\cos x$ on $[0,1]$.
2. Let $f(x)=3(x+1)\left(x-\frac{1}{2}\right)(x-1)$. Use the Bisection method on the following intervals to find $p_{3}$.
(a) $[-2,1.5]$
(b) $[-1.25,2.5]$
3. Use the Bisection method to find solutions accurate to within $10^{-2}$ for $x^{3}-$ $7 x^{2}+14 x-6=0$ on each interval.
(a) $[0,1]$
(b) $[1,3.2]$
(c) $[3.2,4]$
4. Use the Bisection method to find solutions accurate to within $10^{-2}$ for $x^{4}-$ $2 x^{3}-4 x^{2}+4 x+4=0$ on each interval.
(a) $[-2,-1]$
(b) $[0,2]$
(c) $[2,3]$
(d) $[-1,0]$
5. (a) Sketch the graphs of $y=x$ and $y=2 \sin x$.
(b) Use the Bisection method to find an approximation to within $10^{-2}$ to the first positive value of $x$ with $x=2 \sin x$.
6. (a) Sketch the graphs of $y=x$ and $y=\tan x$.
(b) Use the Bisection method to find an approximation to within $10^{-2}$ to the first positive value of $x$ with $x=\tan x$.
7. Let $f(x)=(x+2)(x+1) x(x-1)^{3}(x-2)$. To which zero of $f$ does the Bisection method converge for the following intervals?
(a) $[-3,2.5]$
(b) $[-2.5,3]$
(c) $[-1.75,1.5]$
(d) $[-1.5,1.75]$
8. Let $f(x)=(x+2)(x+1)^{2} x(x-1)^{3}(x-2)$. To which zero of $f$ does the Bisection method converge for the following intervals?
(a) $[-1.5,2.5]$
(b) $[-0.5,2.4]$
(c) $[-0.5,3]$
(d) $[-3,-0.5]$
9. Use the Bisection method to find an approximation to $\sqrt{3}$ correct to within $10^{-4}$. [Hint: Consider $f(x)=x^{2}-3$.]
10. Use the Bisection method to find an approximation to $\sqrt[3]{25}$ correct to within $10^{-4}$.
11. Find a bound for the number of Bisection method iterations needed to achieve an approximation with accuracy $10^{-3}$ to the solution of $x^{3}+x-4=0$ lying in the interval $[1,4]$. Find an approximation to the root with this degree of accuracy.
12. Find a bound for the number of Bisection method iterations needed to achieve an approximation with accuracy $10^{-4}$ to the solution of $x^{3}-x-1=0$ lying in the interval $[1,2]$. Find an approximation to the root with this degree of accuracy.
13. The function defined by $f(x)=\sin \pi x$ has zeros at every integer. Show that when $-1<a<0$ and $2<b<3$, the Bisection method converges to
(a) 0 , if $a+b<2$
(b) 2, if $a+b>2$
(c) 1 , if $a+b=2$
