

جامعة الانبار كلية العلوم قسم الرياضيات نظرية البيانات Introduction to Graphs (2)

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# Lecture (2) Graph Theory Introduction to Graphs (2)

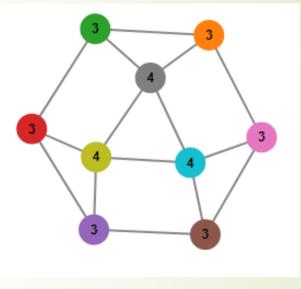
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## Outlines

- ✓ Degree Sequences in Graphs.
- ✓ Graphical Sequence.
- ✓ Havel Hakimi Algorithm.
- ✓ Neighborhoods.
- ✓ H. W.

**Definition:** (Degree Sequence) The degree sequence of a graph of order n is the n –term sequence (usually written in descending order) of the vertex degrees.

Degree Sequence = (4, 4, 4, 3, 3, 3, 3, 3, 3, 3)



**Definition:** (Graphical Sequence) An integer sequence is said to be graphical if it is the degree sequence of some graphs.

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**Example 1:** Is the sequence S = (9,9,8,7,7,6,6,5,5) graphical? Justify your answer.

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**Solution:** The sequence  $S = (a_i)$  is graphical if every element of *S* is the degree of some vertex in a graph. For any graph, we know that  $\sum_{v \in V(G)} d(v) = 2E$ , an even integer. Here,  $\sum a_i = 62$ , an even number. But note that the maximum degree that a vertex can attain in a graph of order *n* is n - 1. If *S* were graphical, the corresponding graph would have been a graph on 9 vertices and have  $\Delta(G) \neq 9$ . Therefore, the given sequence is not graphical.

**Example 2:** Is the sequence S = (9,8,7,6,6,5,5,4,3,3,2,2) graphical? Justify your answer.

**Solution:** The sequence  $S = (a_i)$  is graphical if every element of *S* is the degree of some vertex in a graph. For any graph, we know that  $\sum_{v \in V(G)} d(v) = 2E$ , an even integer. Here, we have  $\sum a_i = 60$ , an even number. Also, note that the all elements in the sequence are less than the number of elements in that sequence. Therefore, the given sequence is graphical.

**Example 3:** Is the sequence S = (5,4,3,3,2,2,2,1,1,1,1) graphical? Justify your answer.

**Solution:** The sequence  $S = (a_i)$  is graphical if every element of *S* is the degree of some vertex in a graph. For any graph, we know that  $\sum_{v \in V(G)} d(v) = 2E$ , an even integer. Here,  $\sum a_i = 25$ , not an even number. Therefore, the given sequence is not graphical.

Havel Hakimi Theorem: The non-negative integer sequence  $D = [d_i]_1^n$  is graphic if and only if  $\hat{D}$  is graphic, where  $\hat{D}$  is the sequence (having n - 1 elements) obtained from D by deleting its largest element  $\Delta$  and subtracting 1 from its  $\Delta$  next largest elements.

# Havel Hakimi Algorithm (HHA)

The Havel Hakimi algorithm gives a systematic approach to answer the question of determining whether it is possible to construct a simple graph from a given degree sequence.

Take as input a degree sequence S and determine if that sequence is graphical That is, can we produce a graph with that degree sequence?

Assume the degree sequence is S

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 $S = d_1, d_2, d_3, \dots, d_n$   $d_i \ge d_{i+1}$ 1. If any  $d_i \ge n$  then fail 2. If there is an odd number of odd degrees then fail 3. If there is a  $d_i < 0$  then fail 4. If all  $d_i = 0$  then report success 5. Reorder S into non - increasing order 6. Let  $k = d_1$ 7. Remove  $d_1$  from S. 8. Subtract 1 from the first k terms remaining of the new sequence 9. Go to step 3 above

### Example 1:

Consider the degree sequence: S = 7, 5, 5, 4, 4, 4, 4, 3

Where n = 8 (no. of vertices)

Step 1. Degree of all vertices is less than n (no.of vertices)

Step 2. Odd number vertices are four.

Step 3. There is no degree less than zero.

Step 4. Remove '7' from the sequence and subtract 1 from the remaining new sequence and arrange again in non-increasing order to get

S = 4, 4, 3, 3, 3, 3, 2

Step 5. Now remove the first '4' from the sequence and subtract 1 from the remaining new sequence to get:

S = 3, 2, 2, 2, 3, 2

rearrange in non-increasing order to get:

S = 3, 3, 2, 2, 2, 2

Repeating the above step we get following degree sequences:

S = 2, 2, 2, 1, 1

S = 1, 1, 1, 1

S = 1, 1, 0

S = 0, 0

**Step 6.** Since all the deg remaining in the sequence is zero, the given sequence is graphical (or in other words, it is possible to construct a simple graph from the given degree sequence).

S = 4, 3, 3, 3, 1

Where n = 5 (no. of vertices)

Step 1. Degree of all vertices is less than n

Step 2. Odd number vertices are four.

Step 3. There is no degree less than zero.

Step 4. Remove '4' from the sequence and subtracting 1 from the remaining new sequence and arrange again in non-increasing order we get

(no.of vertices)

S = 2,2,2,0

Step 5. Again Remove '2 ' from the sequence and subtracting 1 from the remaining new sequence and arrange in non-increasing order we get

S= 1,1,0

Repeating the above step

S= 0,0

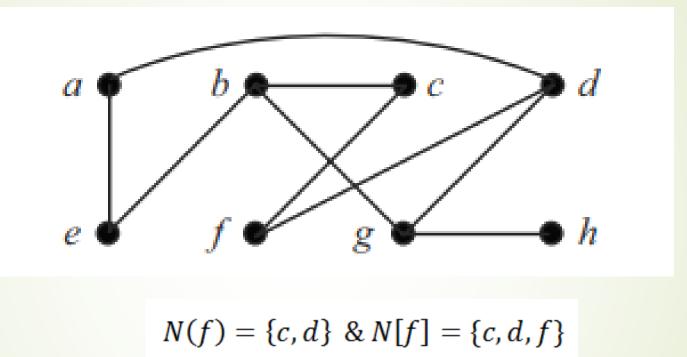
Step 6. Since all the deg remaining in the sequence is zero, the given sequence is graphical.

## 9 Neighbourhoods

**Definition:** (Neighbourhood of a Vertex) The neighbourhood (or open neighbourhood) of a vertex v, denoted by N(v), is the set of vertices adjacent to v. That is,  $N(v) = \{x \in V : vx \in E\}$ . The closed neighbourhood of a vertex v, denoted by N[v], is simply the set  $N(v) \cup \{v\}$ .

Then, for any vertex v in a graph G, we have d(v) = |N(v)|. A special case is a loop that connects a vertex to itself; if such an edge exists, the vertex is said to belong to its own neighbourhood.

Given a set S of vertices, we define the neighbourhood of S, denoted by N(S), to be the union of the neighbourhoods of the vertices in S. Similarly, the closed neighbourhood of S, denoted by N[S], is defined to be  $S \cup N(S)$ .



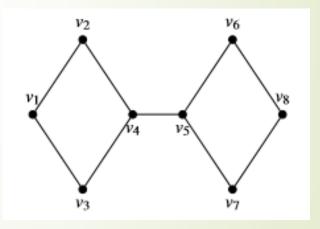
Let  $S = \{a,b,c\}, then N(S) = \{e,d\} \cup \{c,e,g\} \cup \{b,f\} = \{e,d,c,g,b,f\}$  $N[S] = \{e,d\} \cup \{c,e,g\} \cup \{b,f\} \cup \{a\} \cup \{b\} \cup \{c\} = \{e,d,c,g,b,f,a\}.$ 

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### **H. W**.

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- 1) Verify whether the integer sequences (7,6,5,4,3,3,2) and (6,6,5,4,3,3,1) are graphical. (Hint: Use Havel Hakimi Algorithm)
- 2) For the following graph *G*, find:  $\delta(G)$ ,  $\Delta(G)$ ,  $N[v_5]$  and degree sequence.





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