

Trees (3)

Insert Node and Delete Node

الدكتور
اثير العاني

Insert Node

Given a sorted list:

5 8 9 13 21 44 45 46

Insert Node

Given a sorted list:

5 8 9 13 21 44 45 46

If 14 is inserted, the list become:

5 8 9 13 14 21 44 45 46

Insert Node

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If 14 is inserted, the list become:

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Insert Node

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FEATURES OF A SORTED LIST IS PRESERVED

Insert Node

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5 8 9 13 21 44 45 46

If 14 is inserted, the list become:

5 8 9 13 14 21 44 45 46



FEATURES OF A SORTED LIST IS PRESERVED

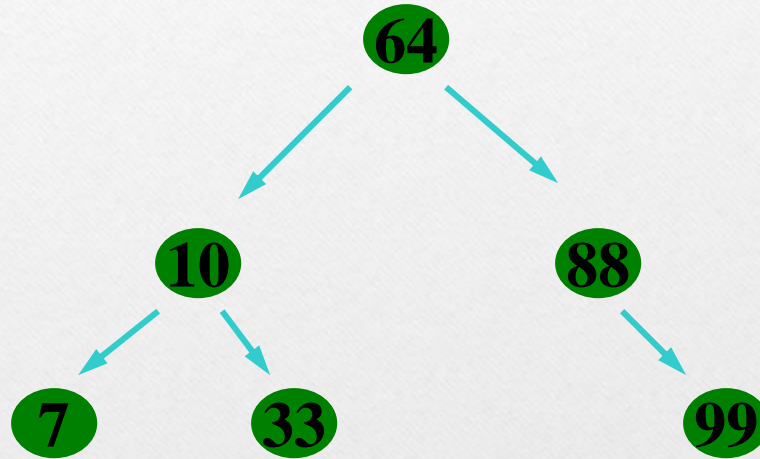
WHAT ARE THE FEATURES OF A SORTED LIST?????

Insert Node

Given a BST:

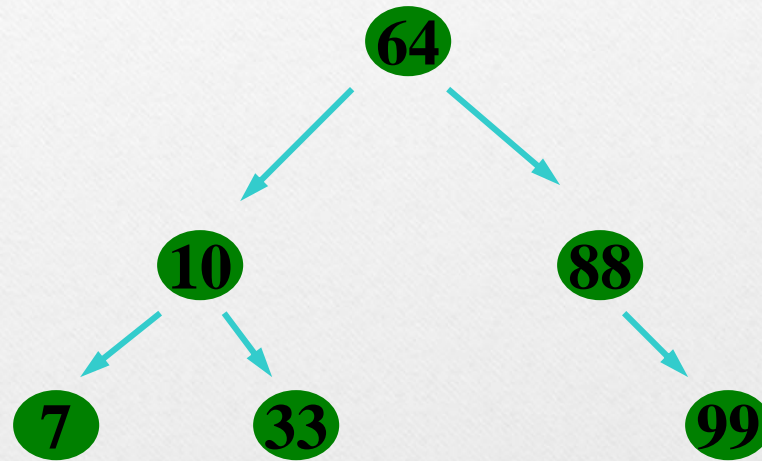
Insert Node

Given a BST:



Insert Node

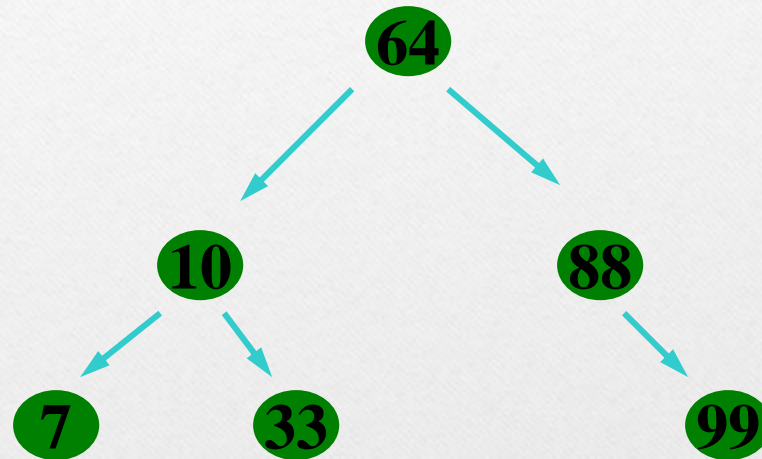
Given a BST:



If 14 to be inserted :

Insert Node

Given a BST:

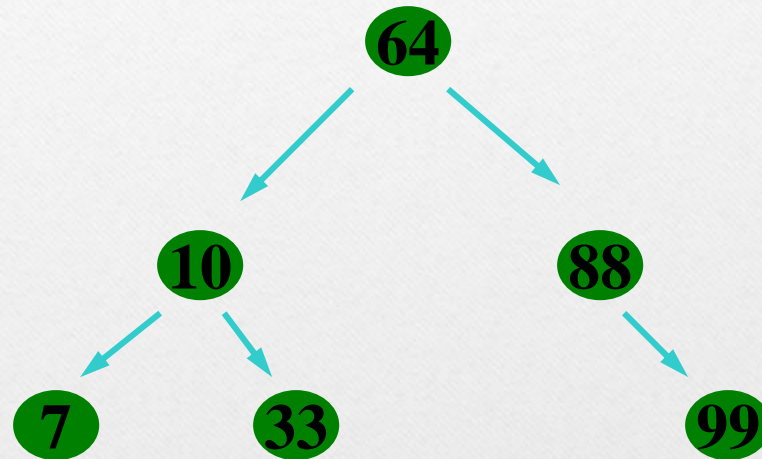


If 14 to be inserted :

How does the new BST look like ?

Insert Node

Given a BST:

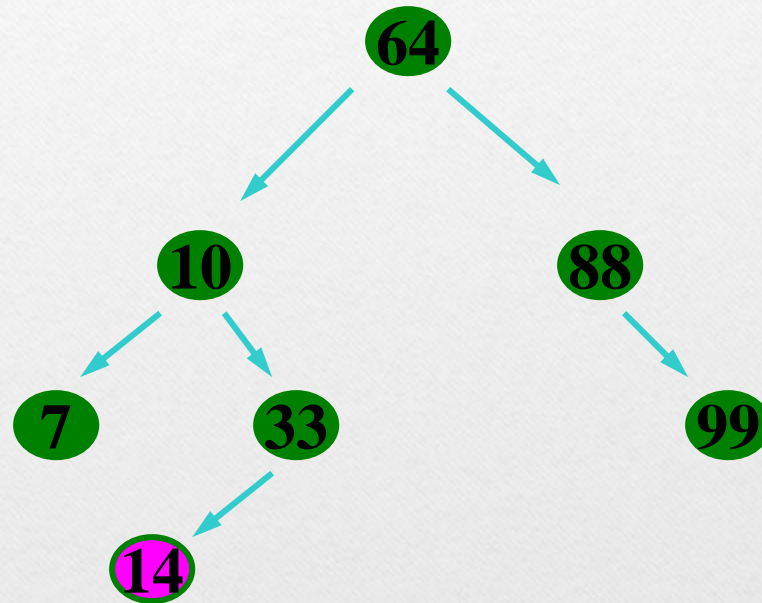


If 14 to be inserted :

**How does the new BST look like ?
How do you do it?**

Insert Node

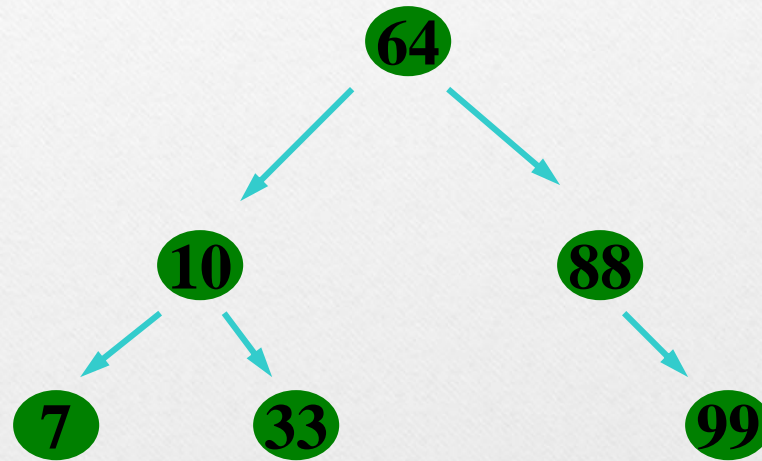
Given a BST:



Are the features of a BST preserved???

Algorithm to insert 14:

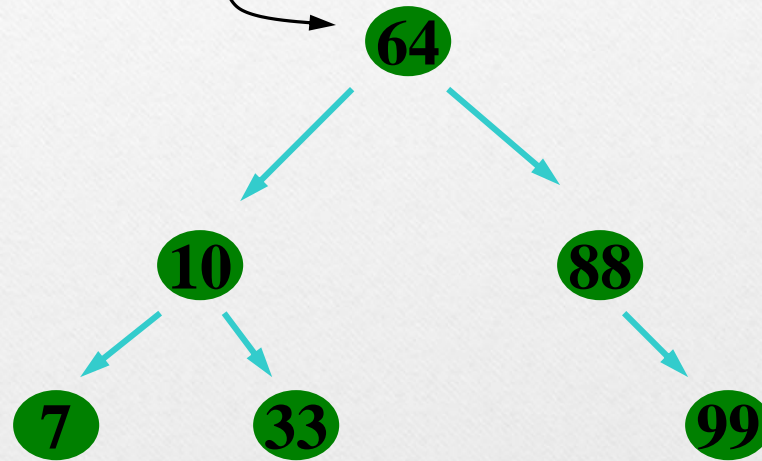
Insert Node



Algorithm to insert 14:

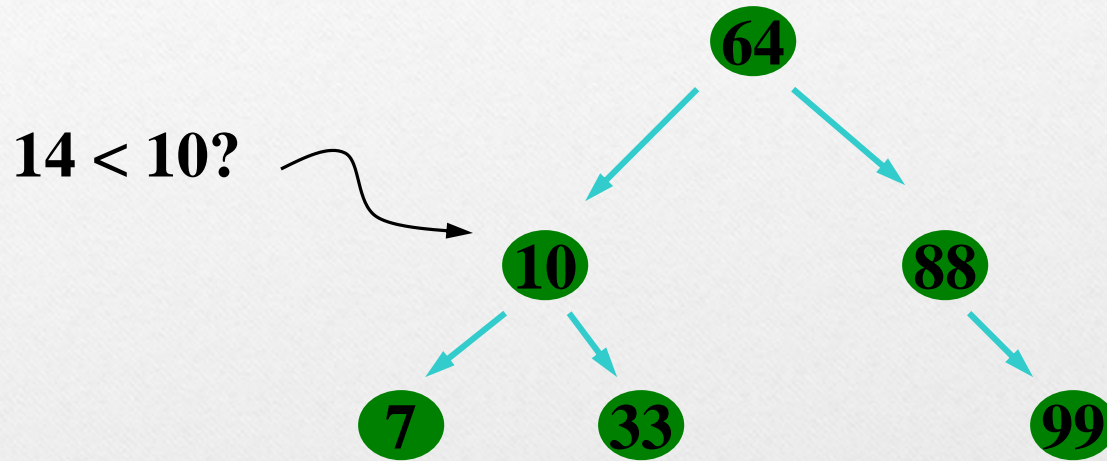
Insert Node

$14 < 64?$



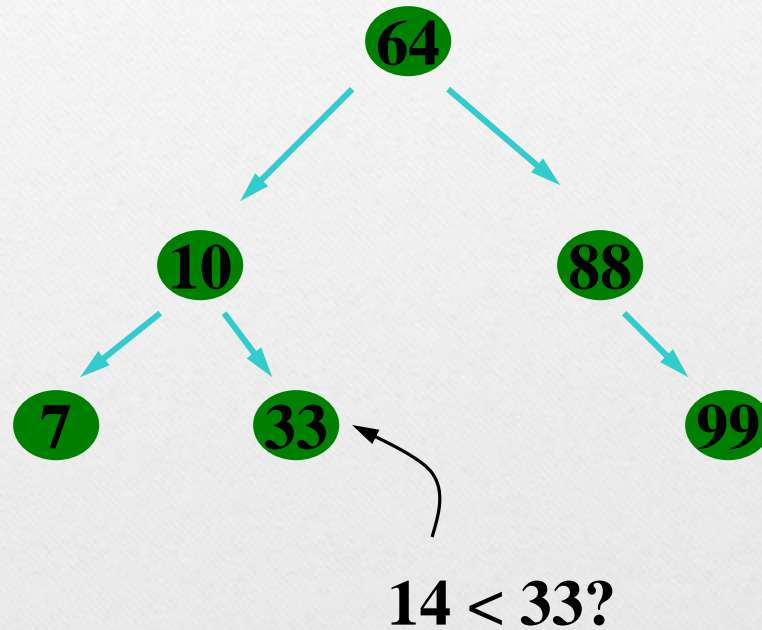
Algorithm to insert 14:

Insert Node



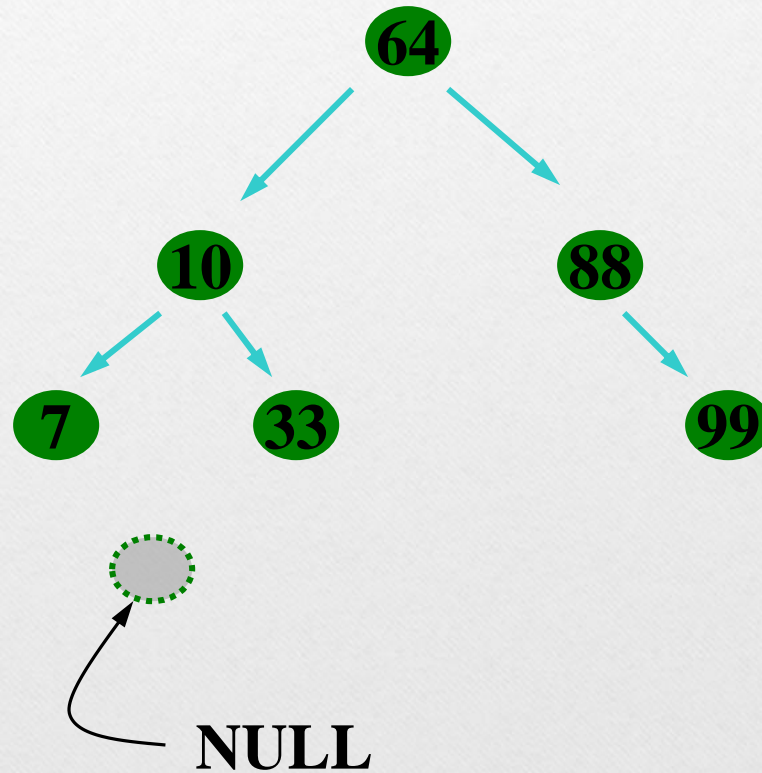
Algorithm to insert 14:

Insert Node



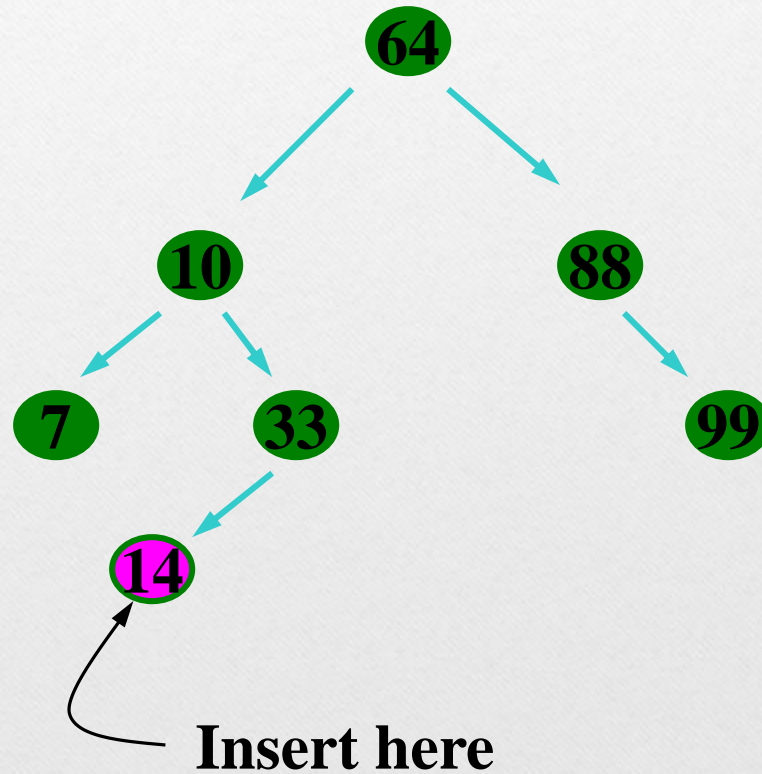
Algorithm to insert 14:

Insert Node



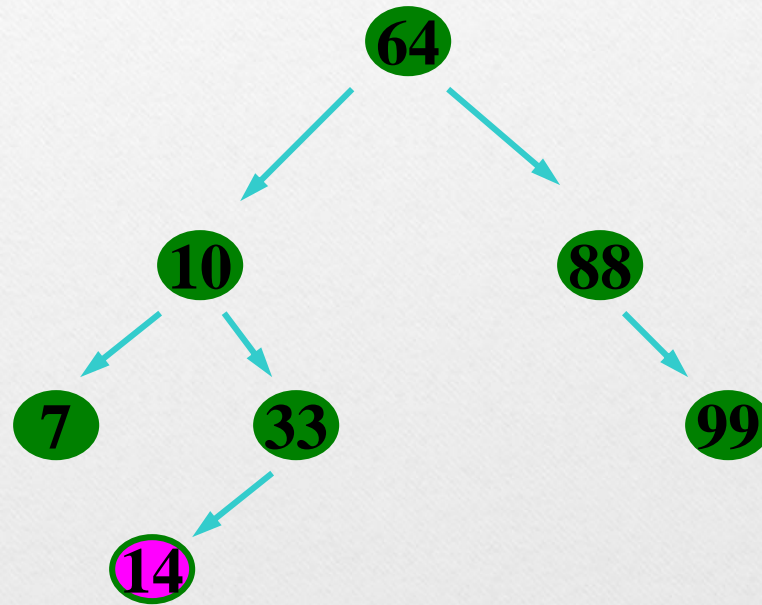
Algorithm to insert 14:

Insert Node



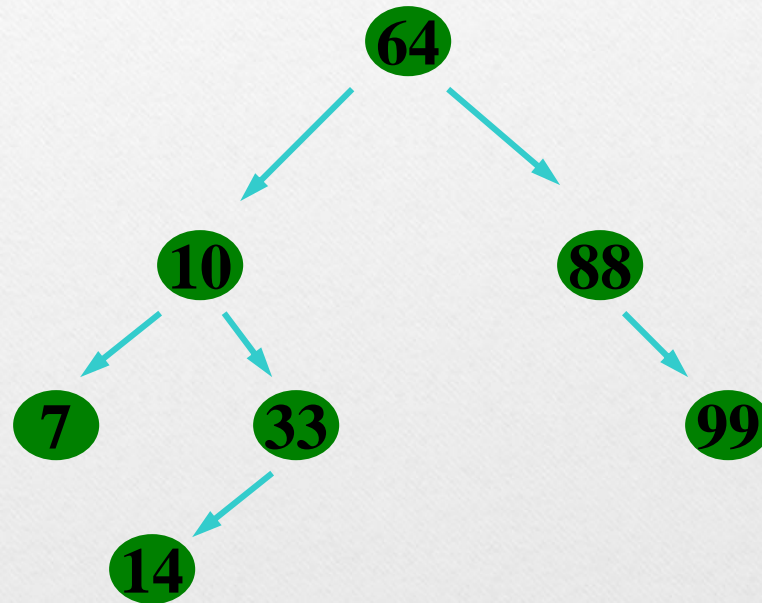
Algorithm to insert 14:

Insert Node



Try insert 99:

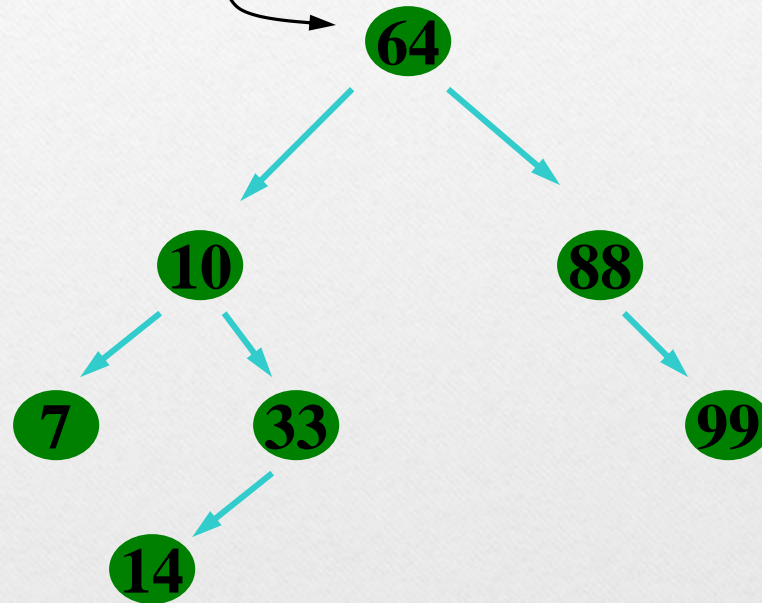
Insert Node



Try insert 99:

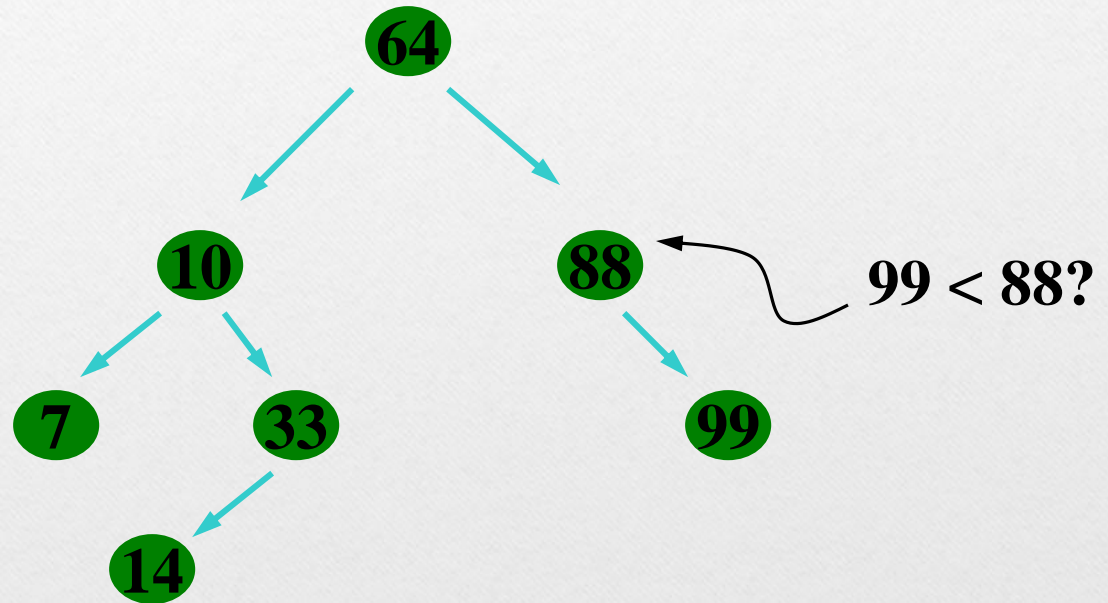
Insert Node

$99 < 64?$



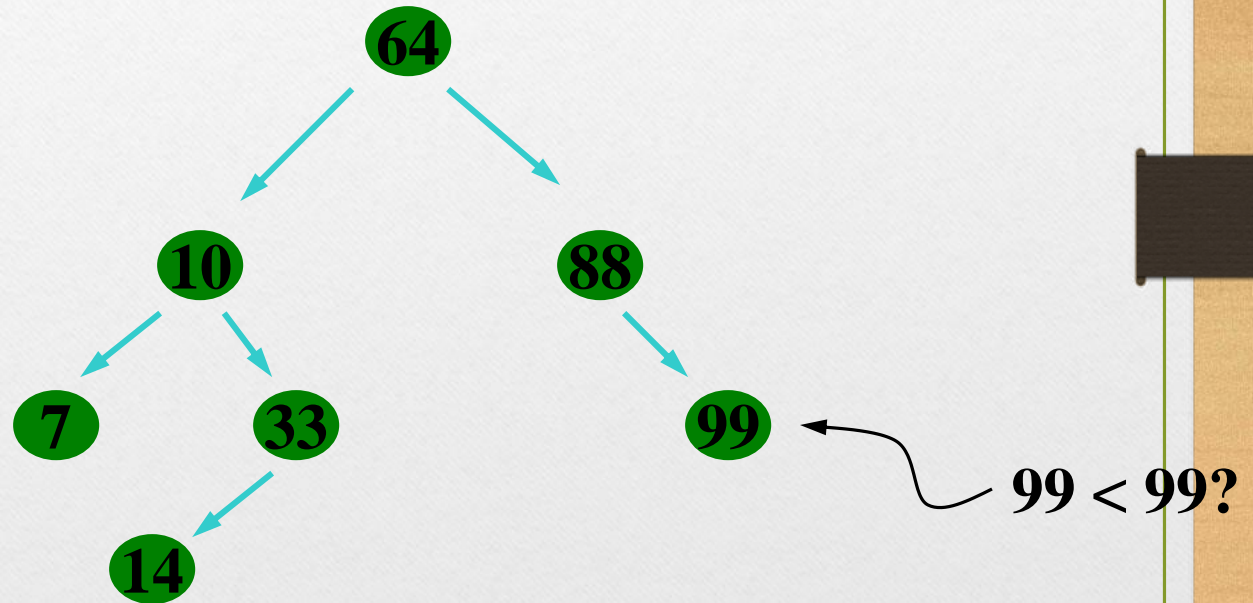
Try insert 99:

Insert Node



Try insert 99:

Insert Node



Insert Node

**This insert algorithm is
not very effective**

**Try building a BST using the
Same algorithm for the following
Sequence:**

K N G I C U

Insert Node

**This insert algorithm is
not very effective**

**Try building a BST using the
Same algorithm for the following
Sequence:**

K N G I C U

K U C I N G

Insert Node

**This insert algorithm is
not very effective**

**Try building a BST using the
Same algorithm for the following
Sequence:**

**K N G I C U
K U C I N G
C G I K N U**

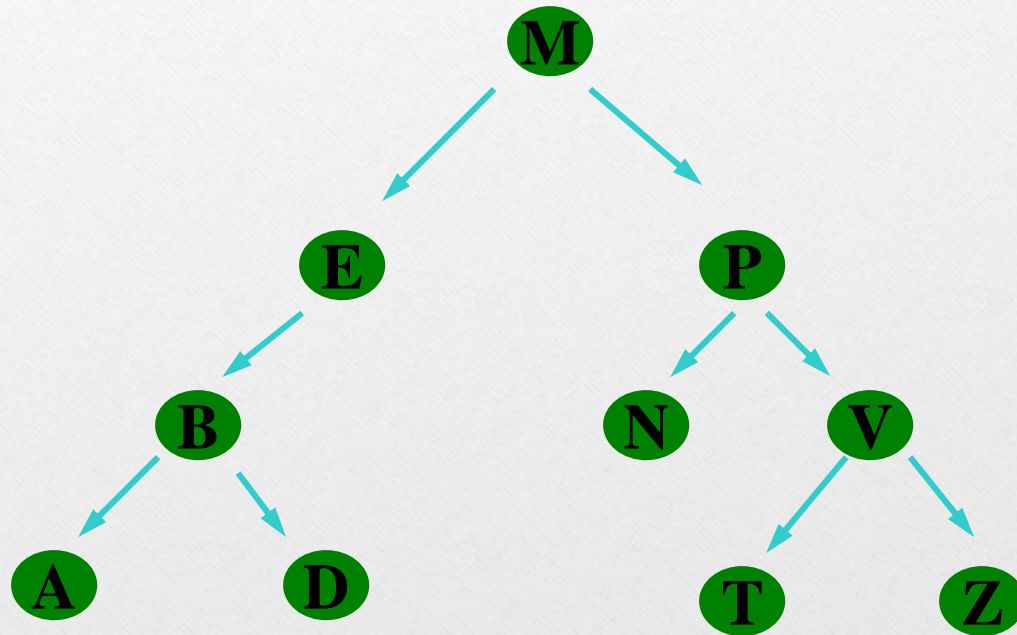
Delete Node

3 cases node deletion :

- The node to be deleted is a leaf
- The node to be deleted has 1 child
- The node to be deleted has 2 children

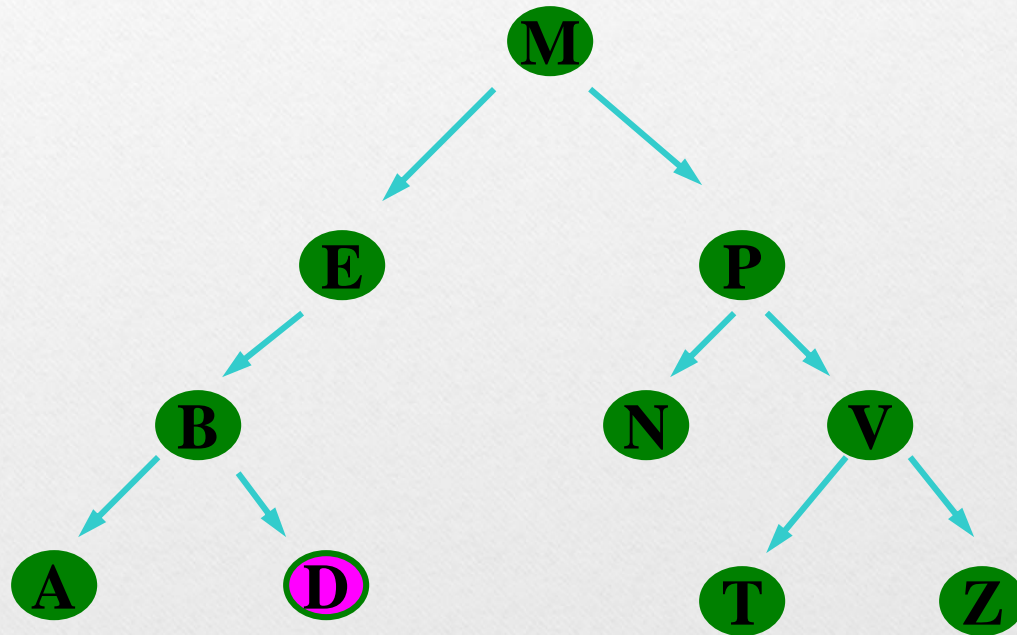
Delete Node: First Case

To delete D:



Delete Node: First Case

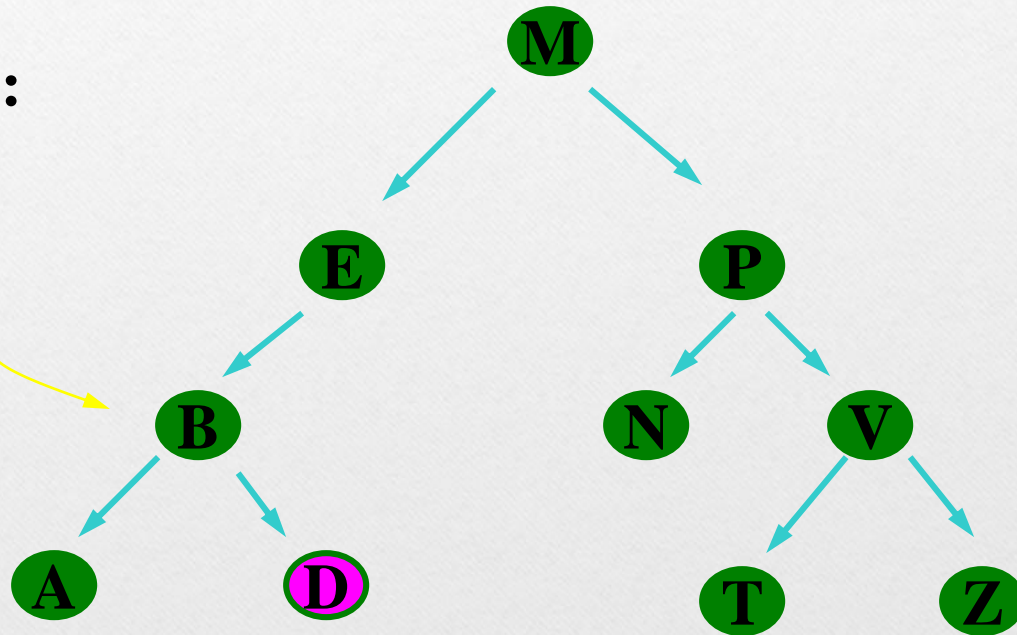
To delete D:



Delete Node: First Case

To delete D:

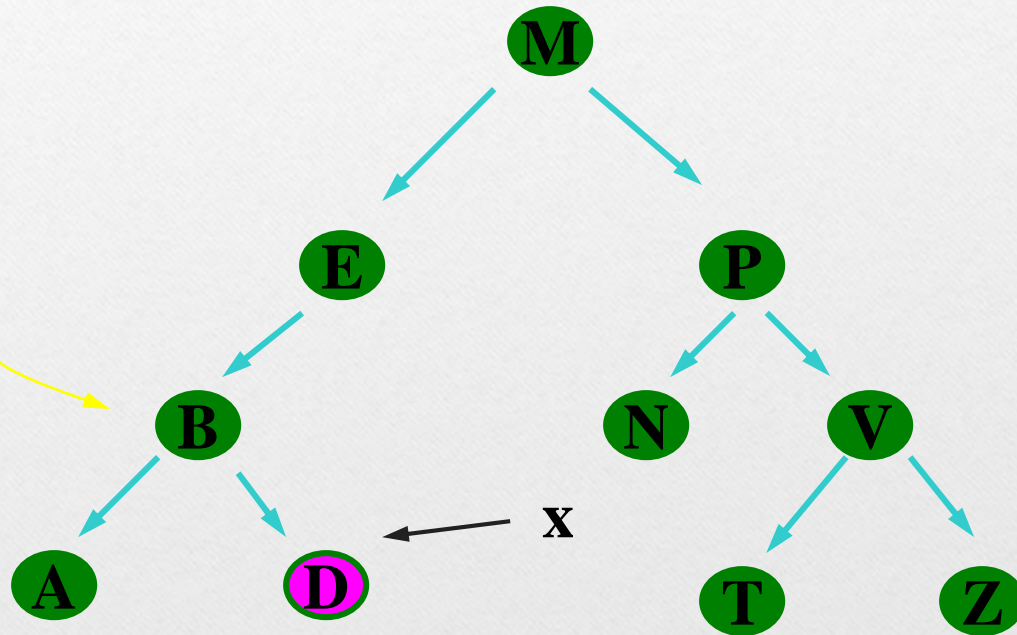
parent



Delete Node: First Case

To delete D

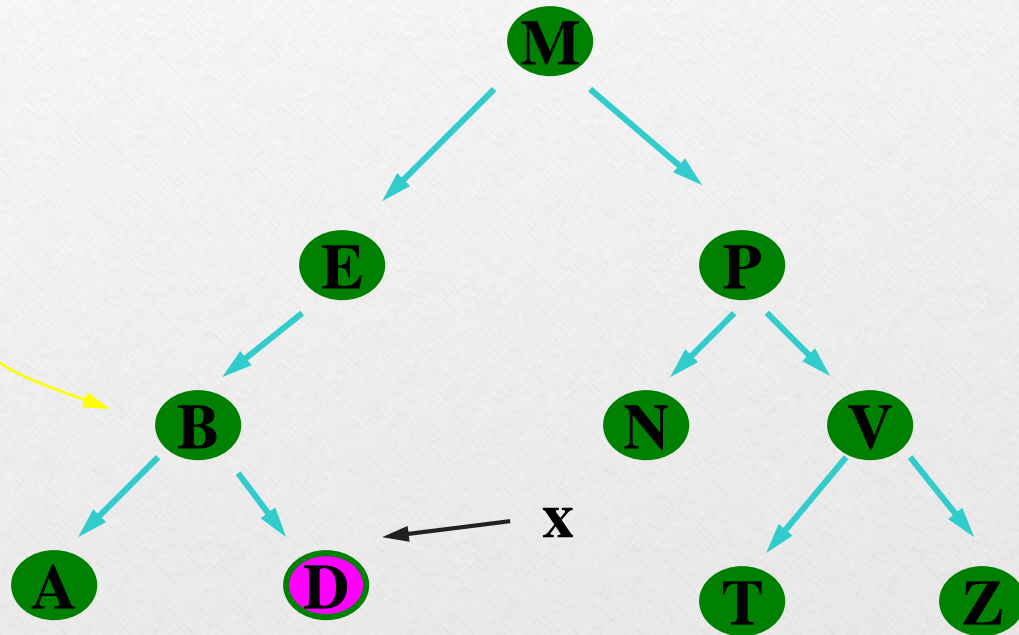
parent



Delete Node: First Case

To delete D

parent

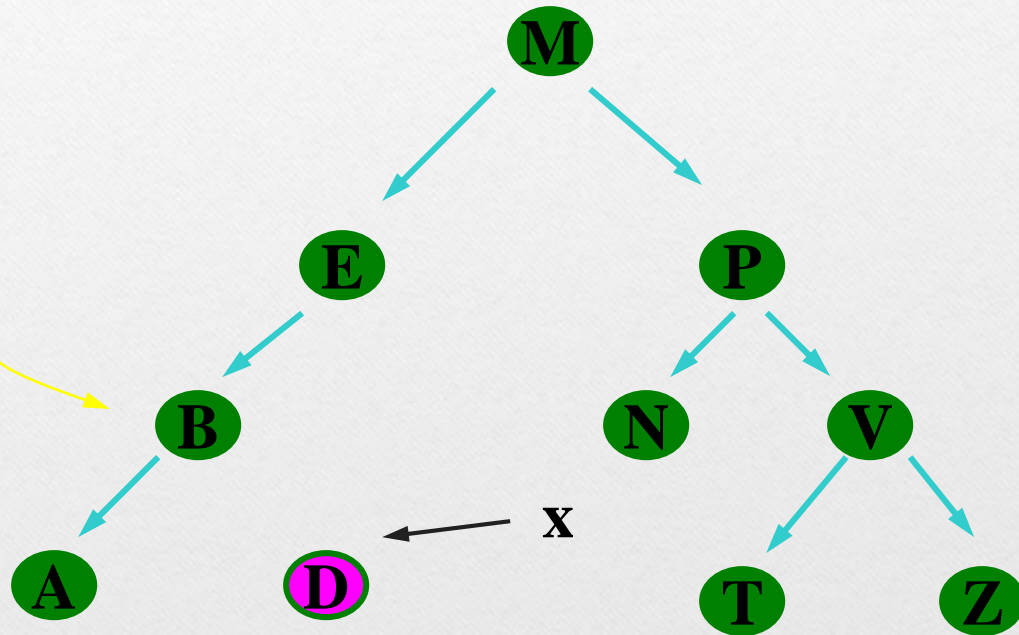


Set the right child of parent as null

Delete Node: First Case

To delete D

parent

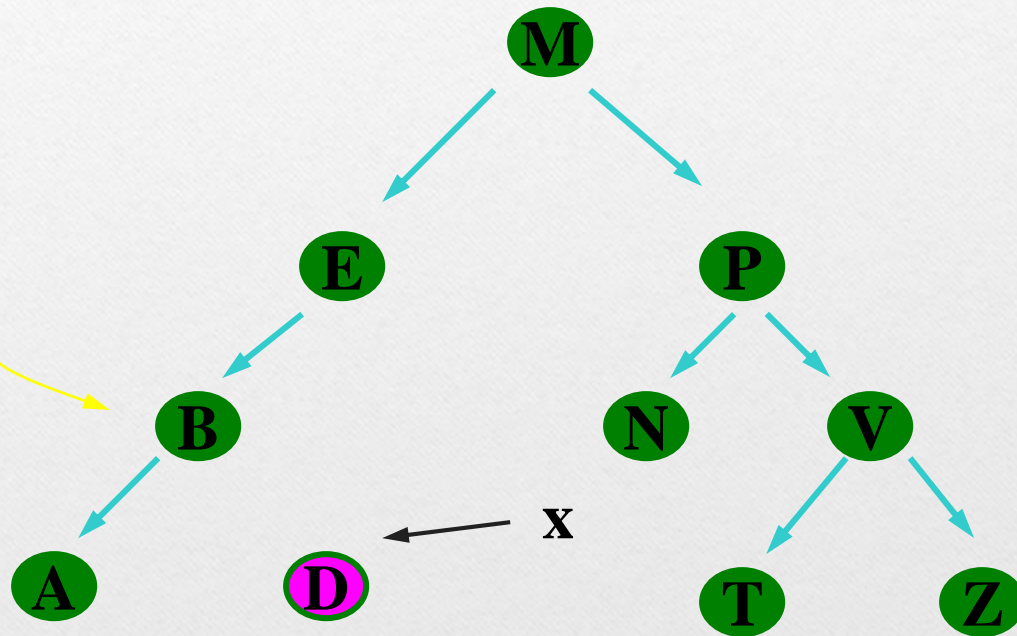


Set the right child of parent as null

Delete Node: First Case

To delete D

parent



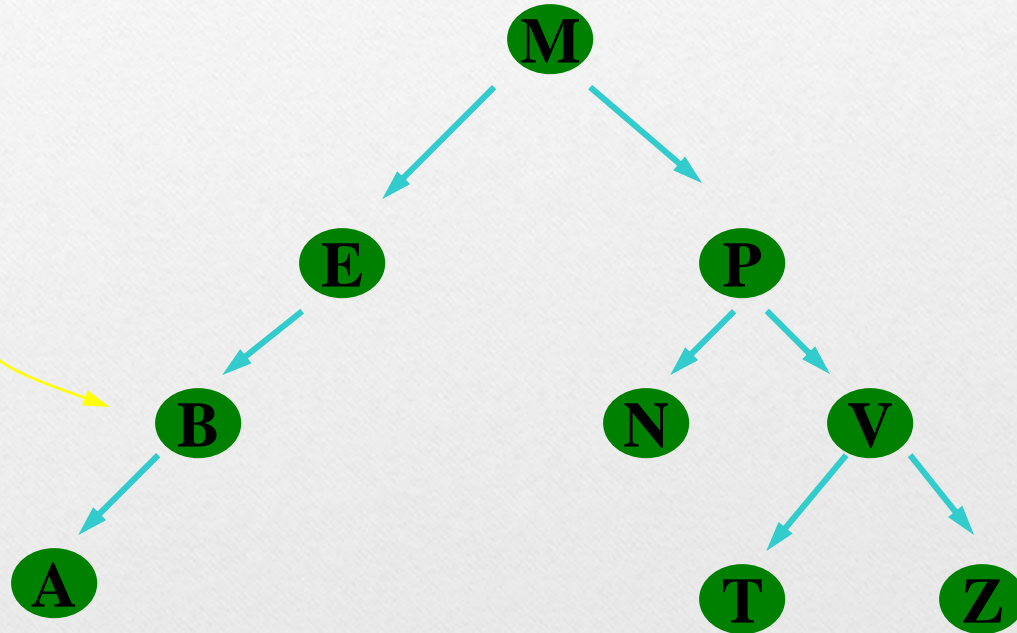
Set the right child of parent as null

Free x

Delete Node: First Case

To delete D

parent

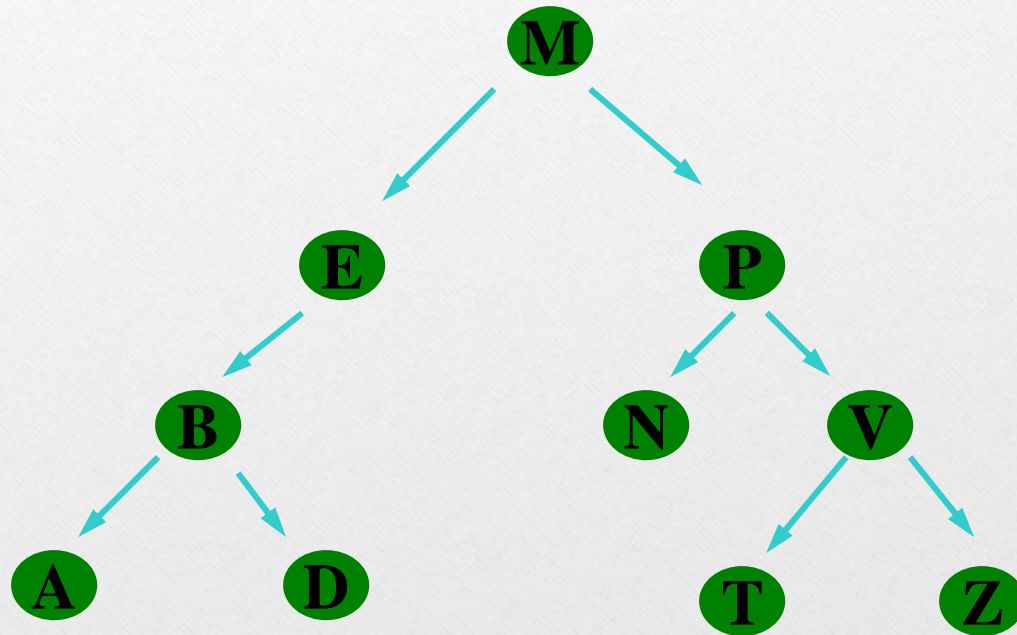


Set the right child of parent as null

Free x

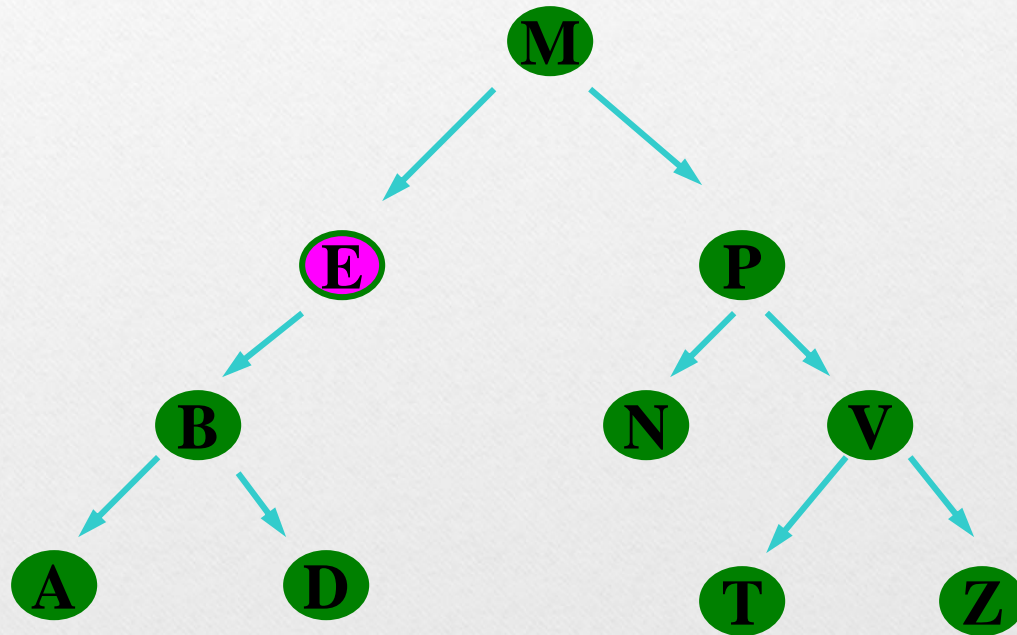
Delete Node: Second Case

To delete E



Delete Node: Second Case

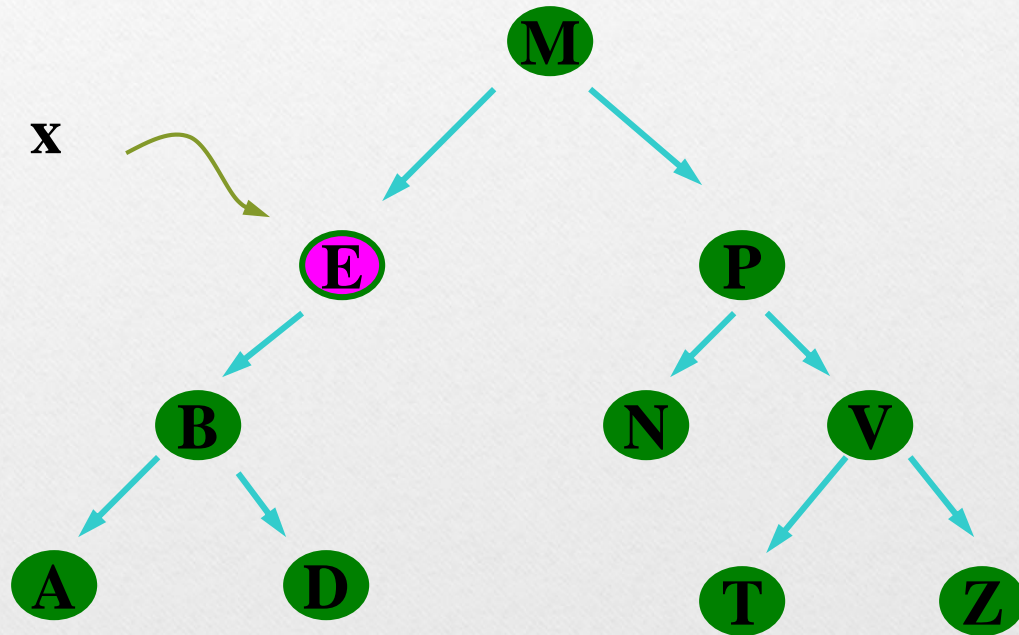
To delete E



Delete Node: Second Case

parent

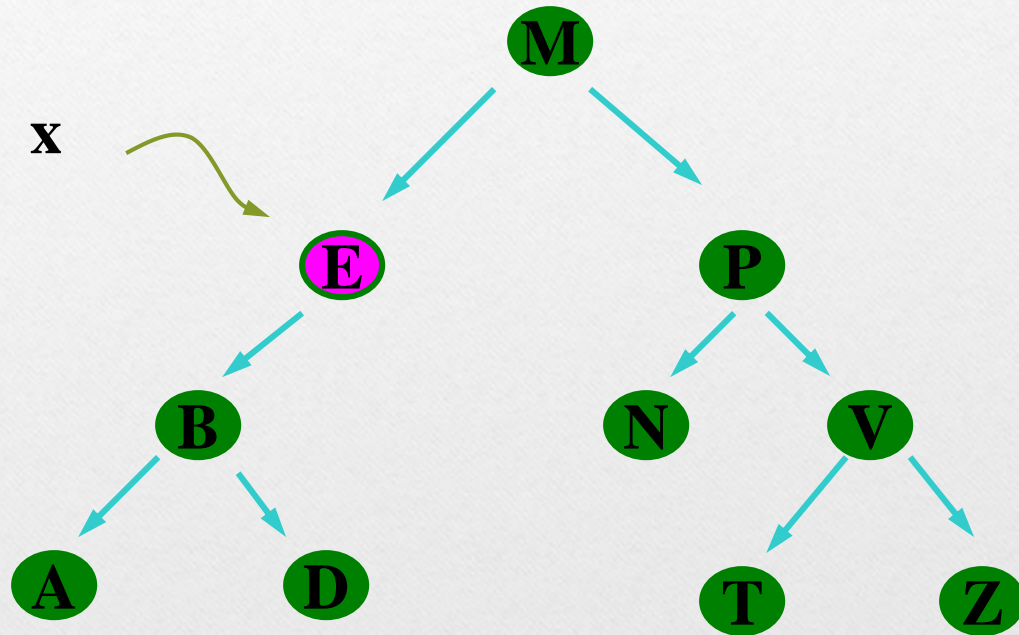
To delete E



Delete Node: Second Case

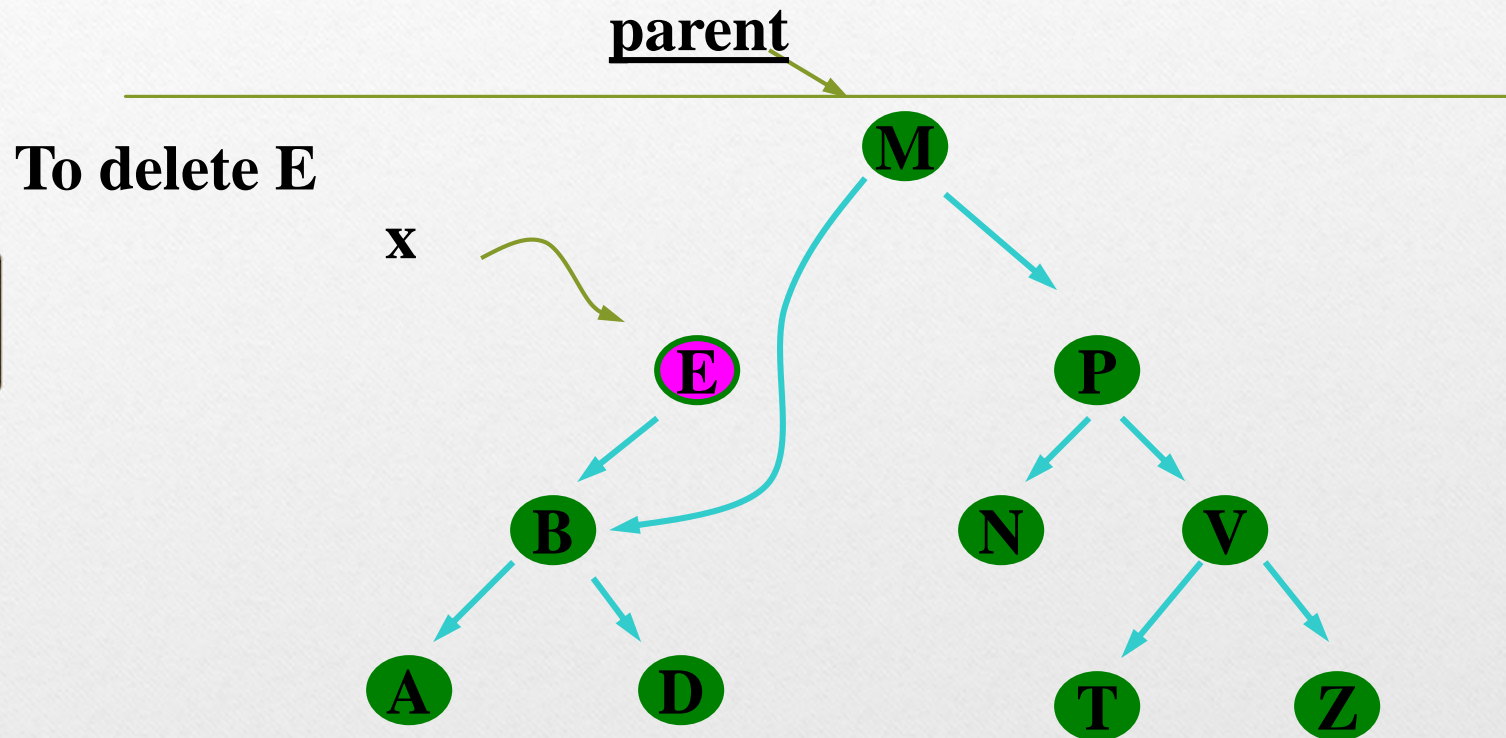
parent

To delete E



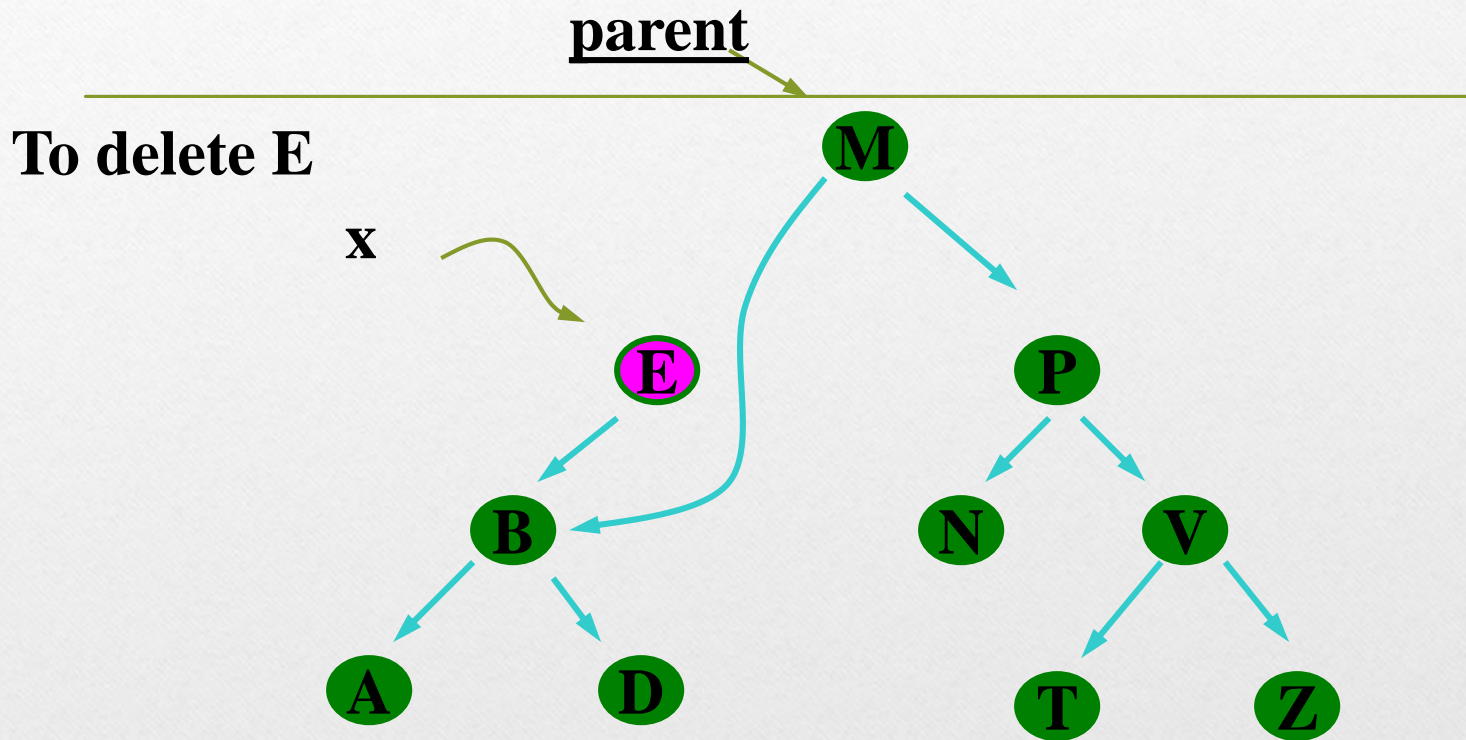
Set the Lchild (@ Rchild) of x as the Lchild (@ Rchild) of parent

Delete Node: Second Case



Set the Lchild (@ Rchild) of x as the Lchild (@ Rchild) of parent

Delete Node: Second Case

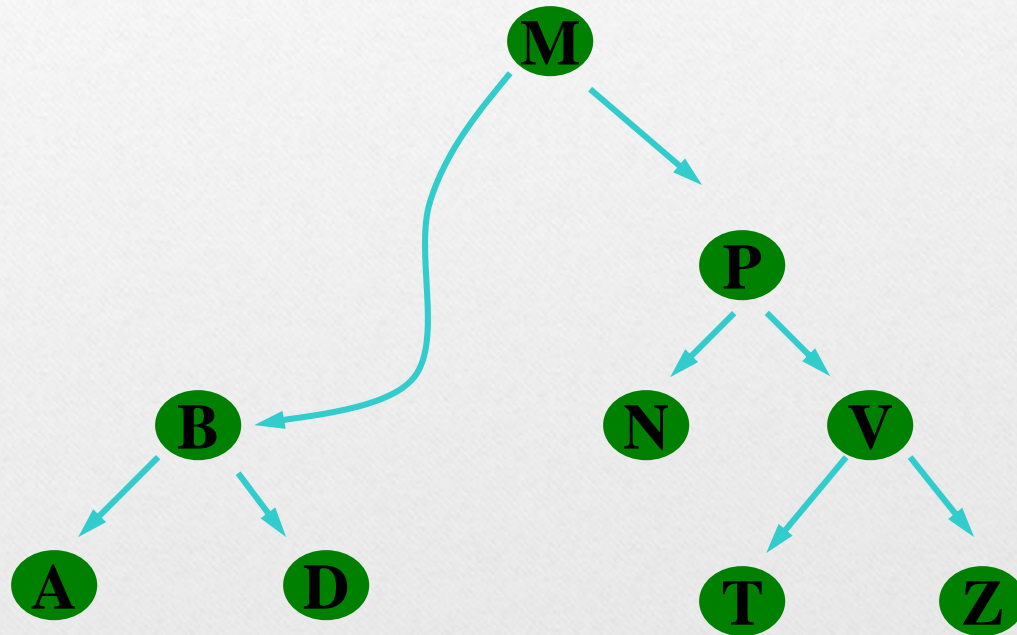


Set the Lchild (@ Rchild) of x as the Lchild (@ Rchild) of parent
Free x

Delete Node: Second Case

parent

To delete E

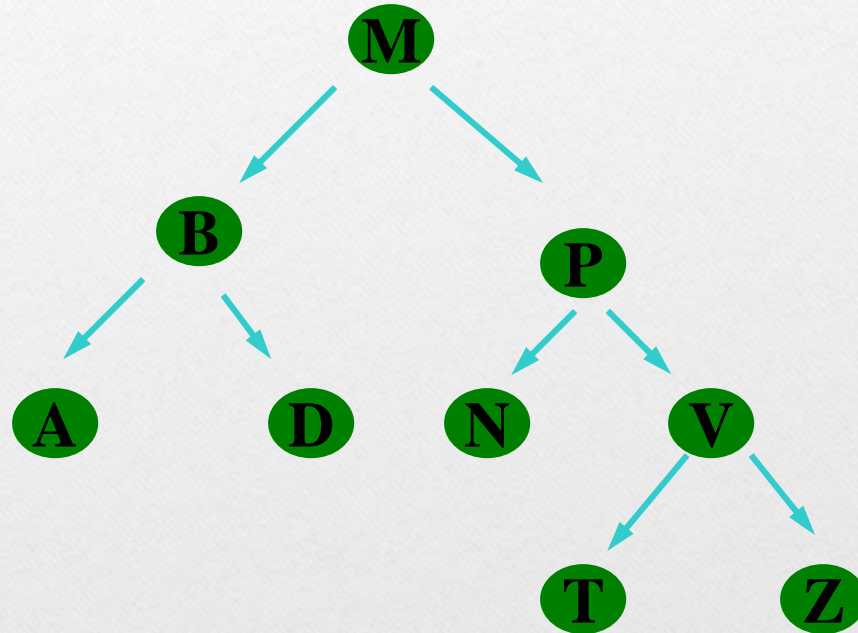


Set the Lchild (@ Rchild) of x as the Lchild (@ Rchild) of parent
Free x

Delete Node: Second Case

parent

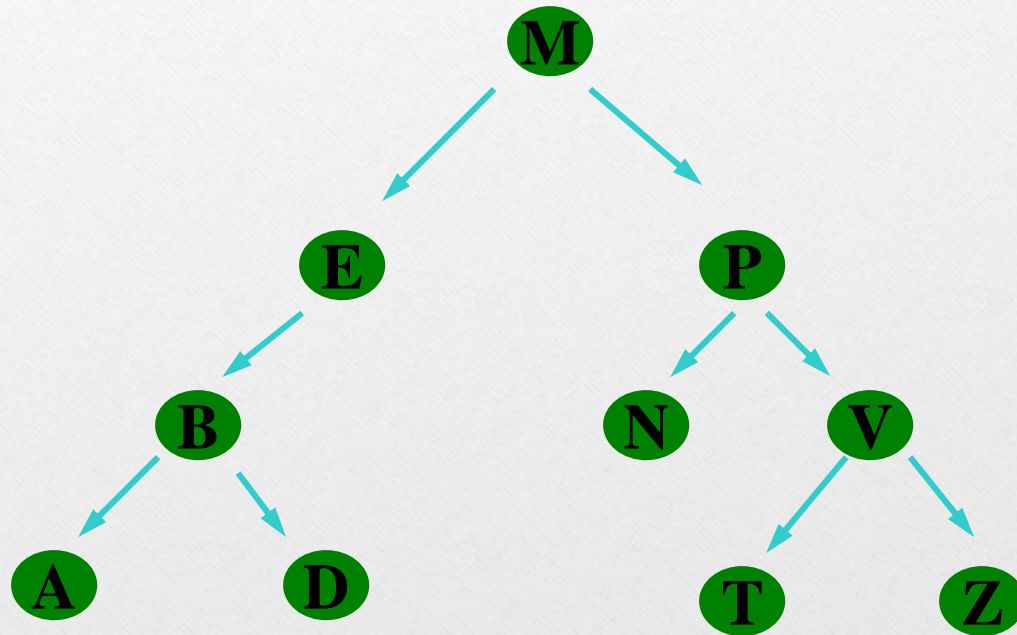
To delete E



Set the Lchild (@ Rchild) of x as the Lchild (@ Rchild) of parent
Free x

Delete Node: Third Case

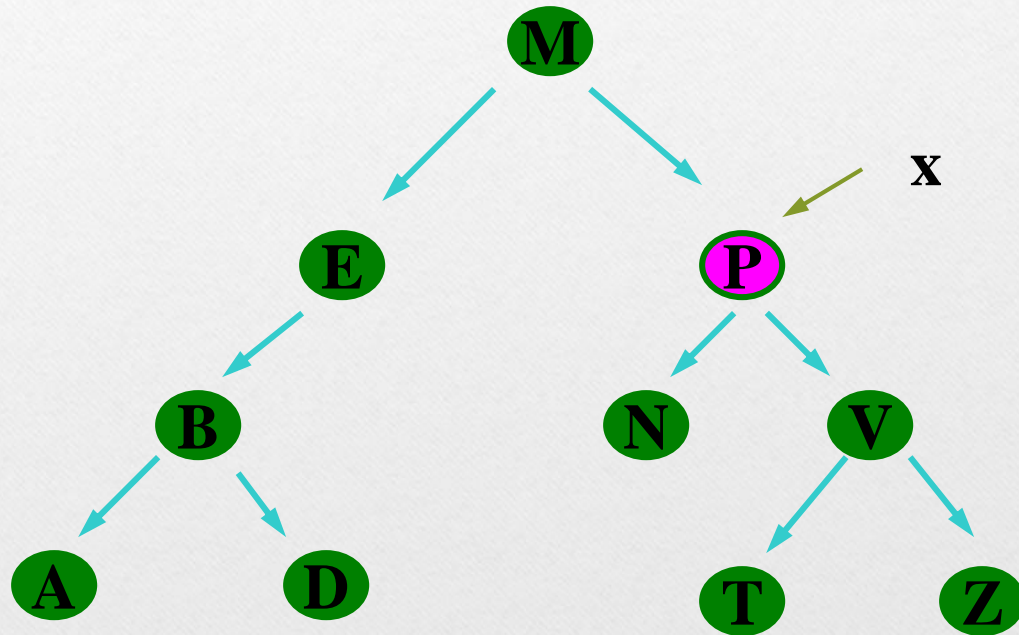
To delete P



Delete Node: Third Case

parent

To delete P

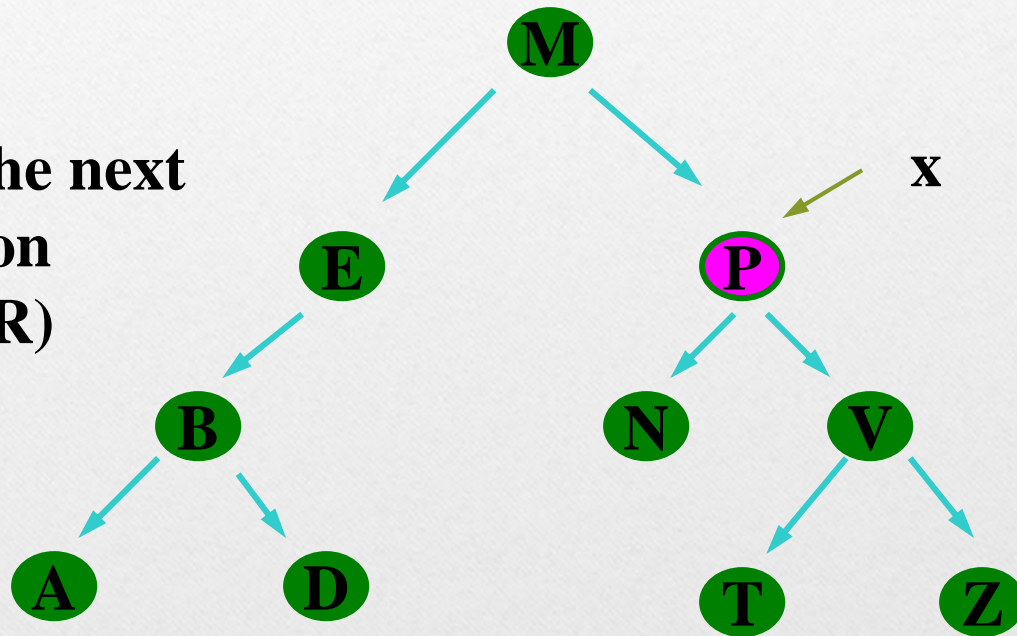


Delete Node: Third Case

parent

To delete P

Determine the next node based on Inorder(LNR)



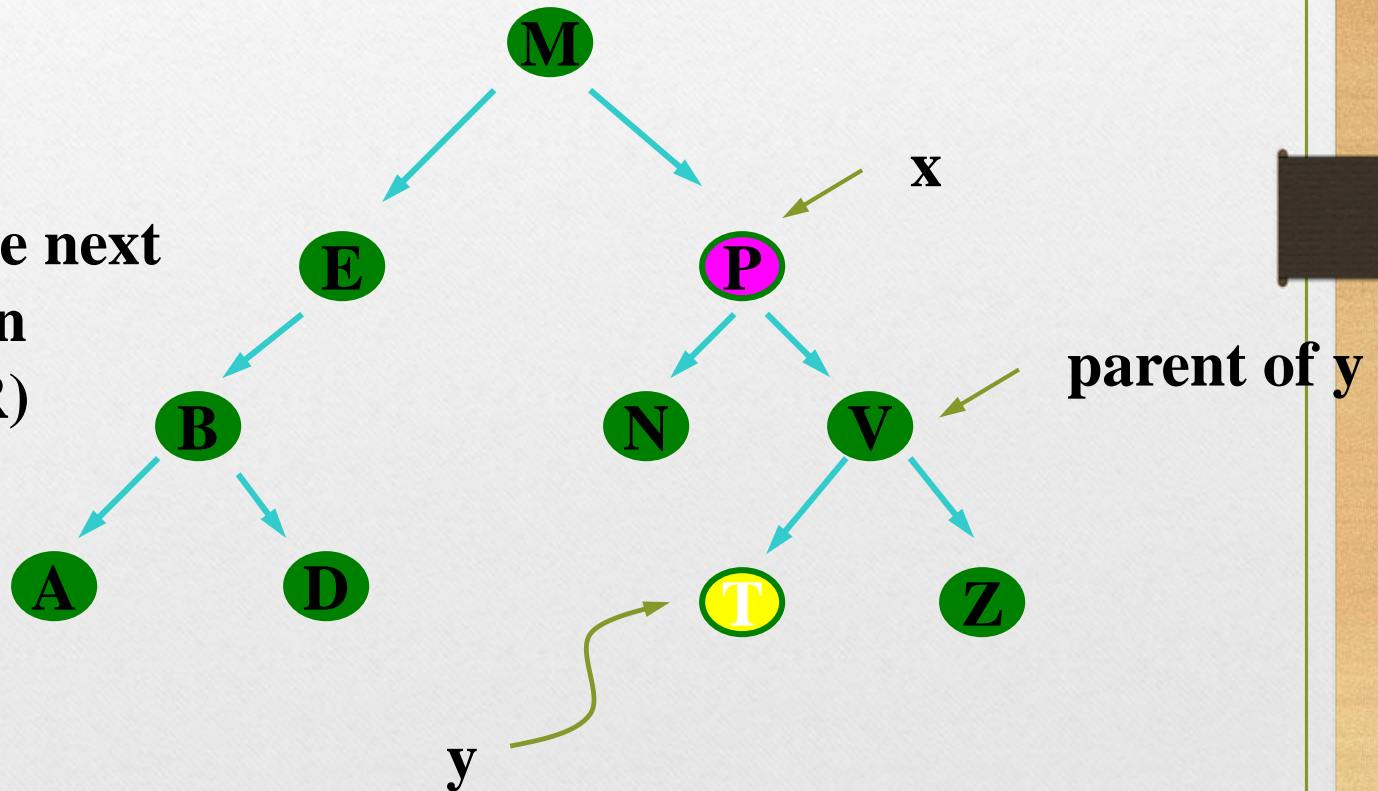
*Usually, the Inorder node has only one child or no children at all

Delete Node: Third Case

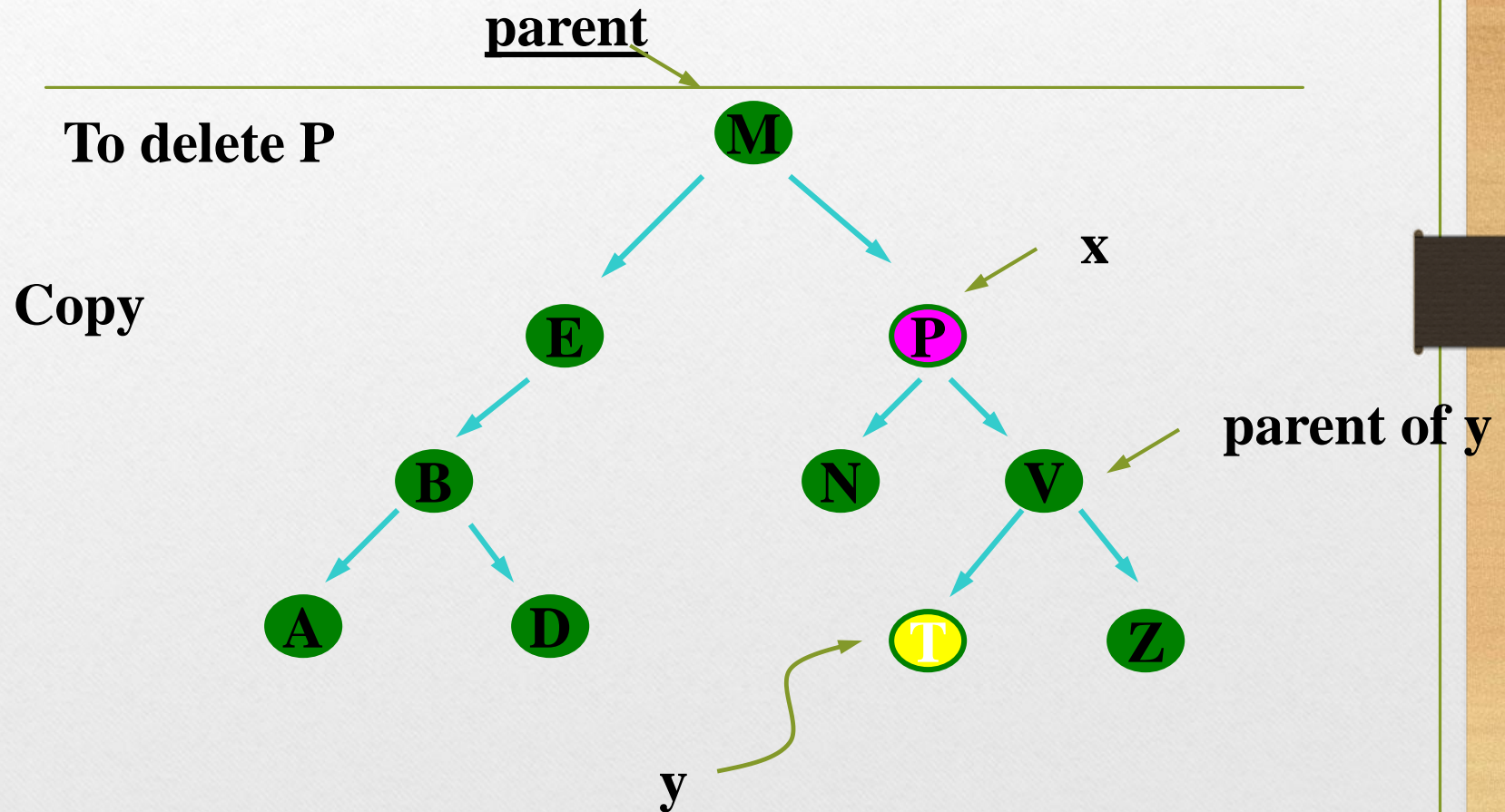
parent

To delete P

Determine the next
node based on
Inorder(LNR)



Delete Node: Third Case



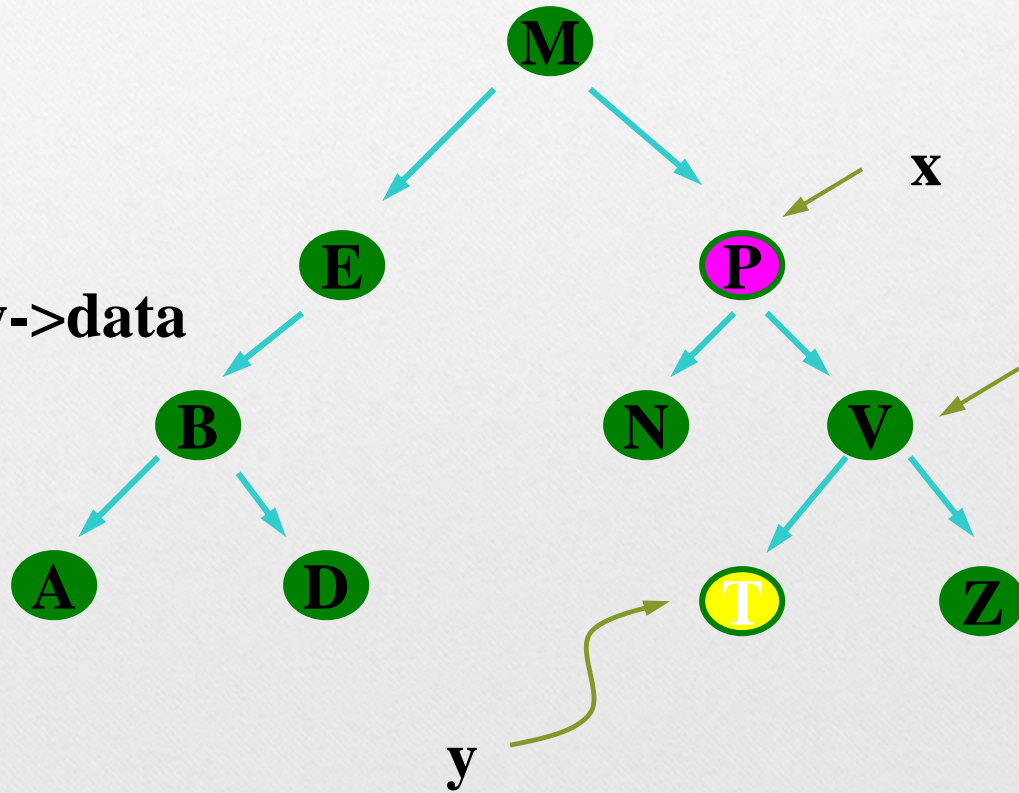
Delete Node: Third Case

parent

To delete P

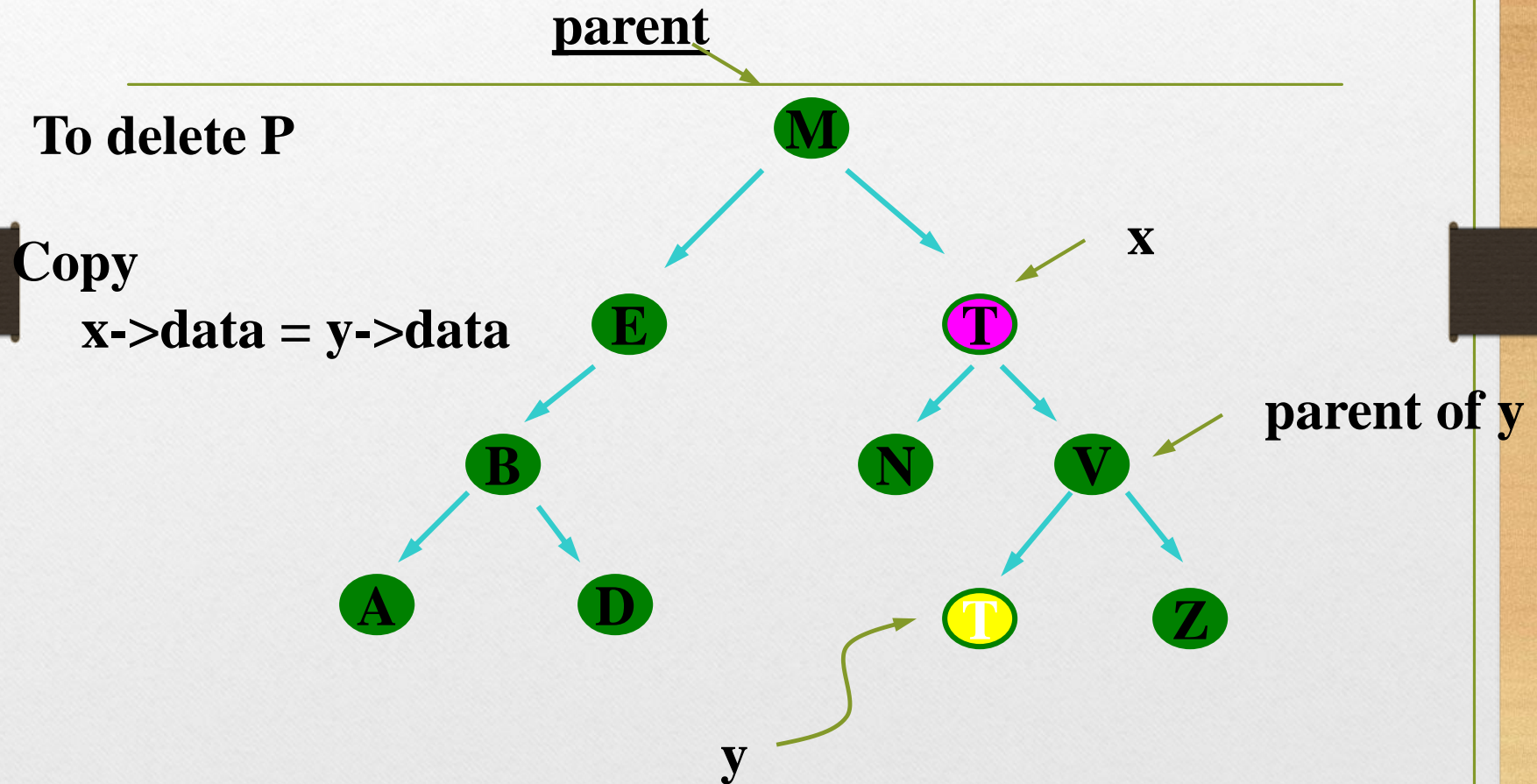
Copy

$x \rightarrow \text{data} = y \rightarrow \text{data}$



parent of y

Delete Node: Third Case



Delete Node: Third Case

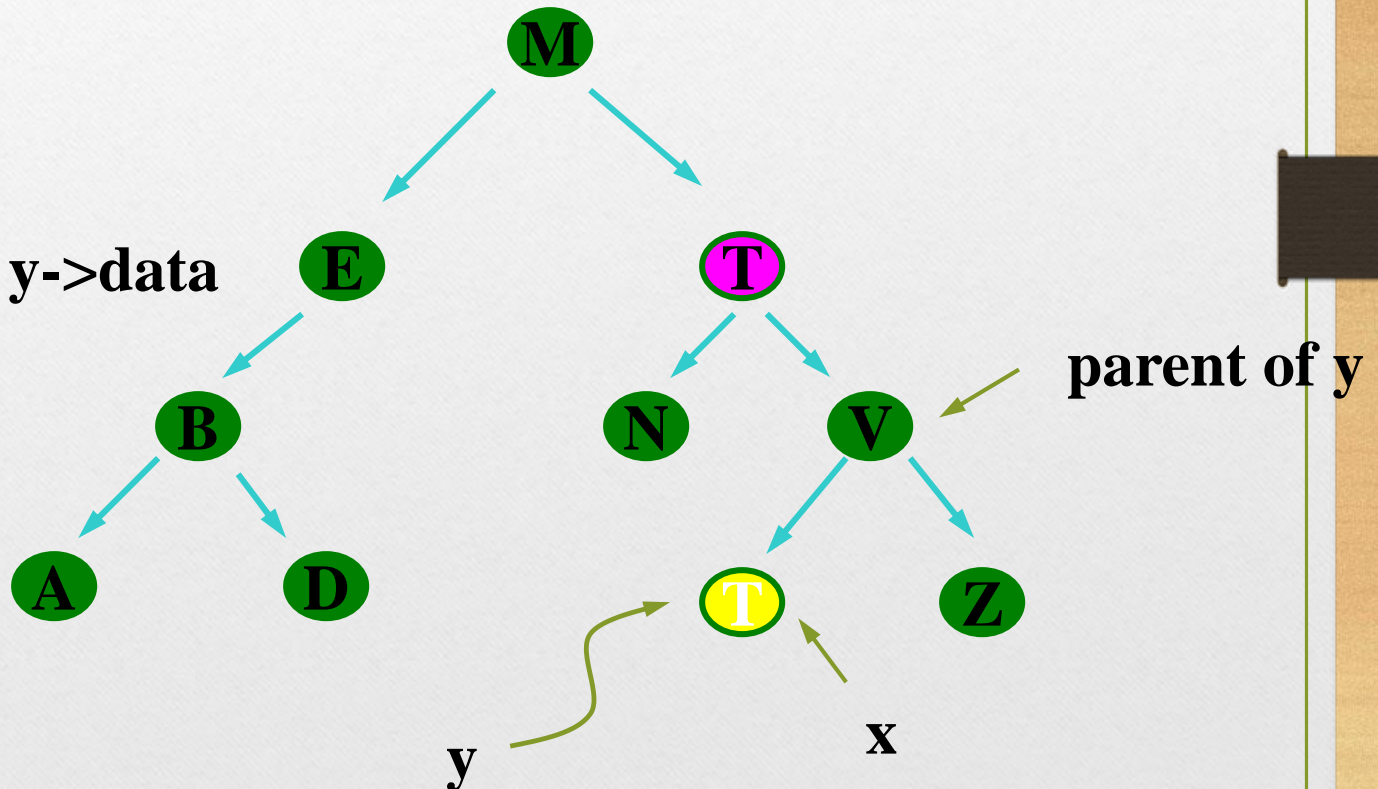
parent

To delete P

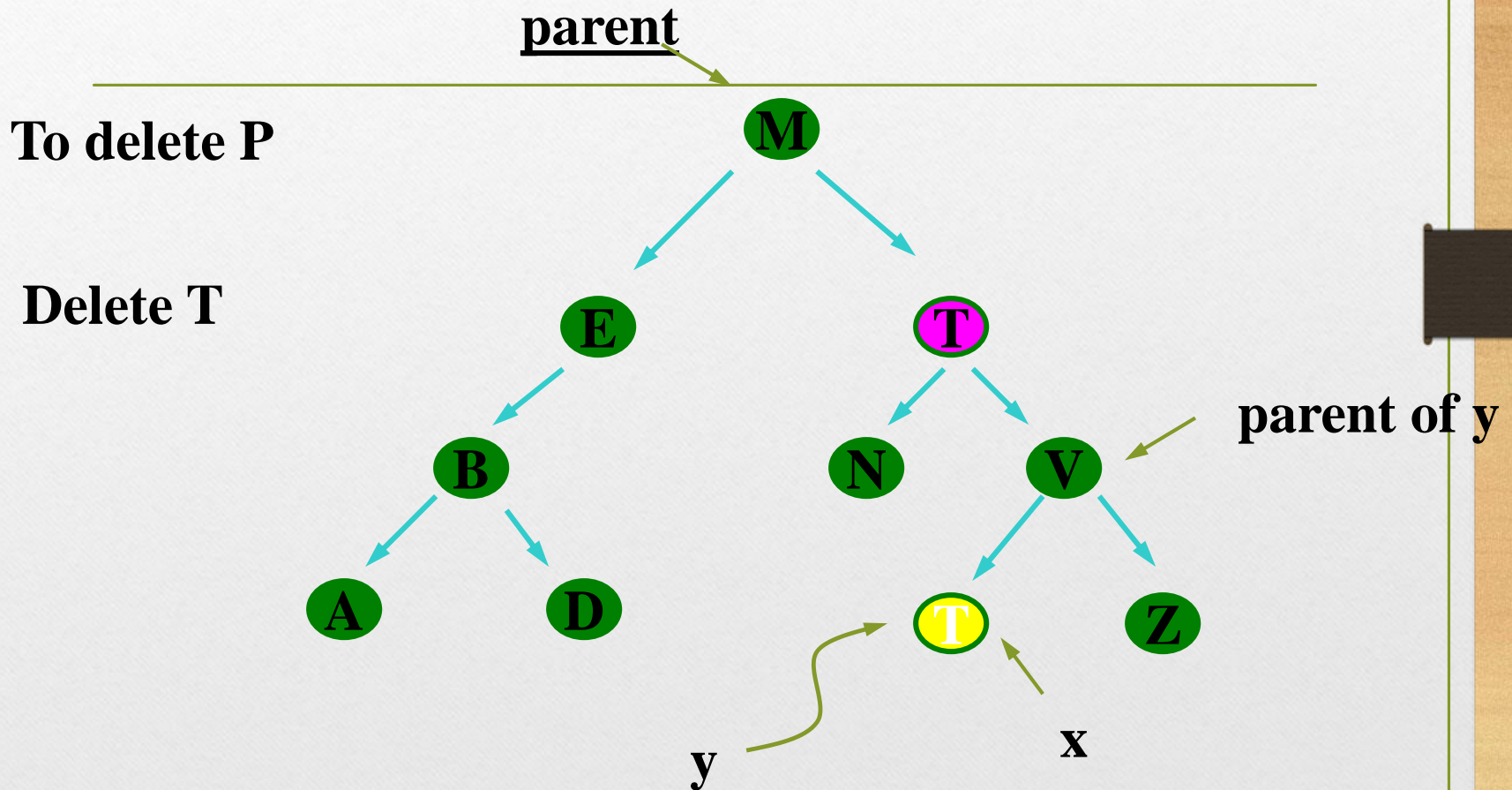
Copy

$x \rightarrow \text{data} = y \rightarrow \text{data}$

$x = y$



Delete Node: Third Case



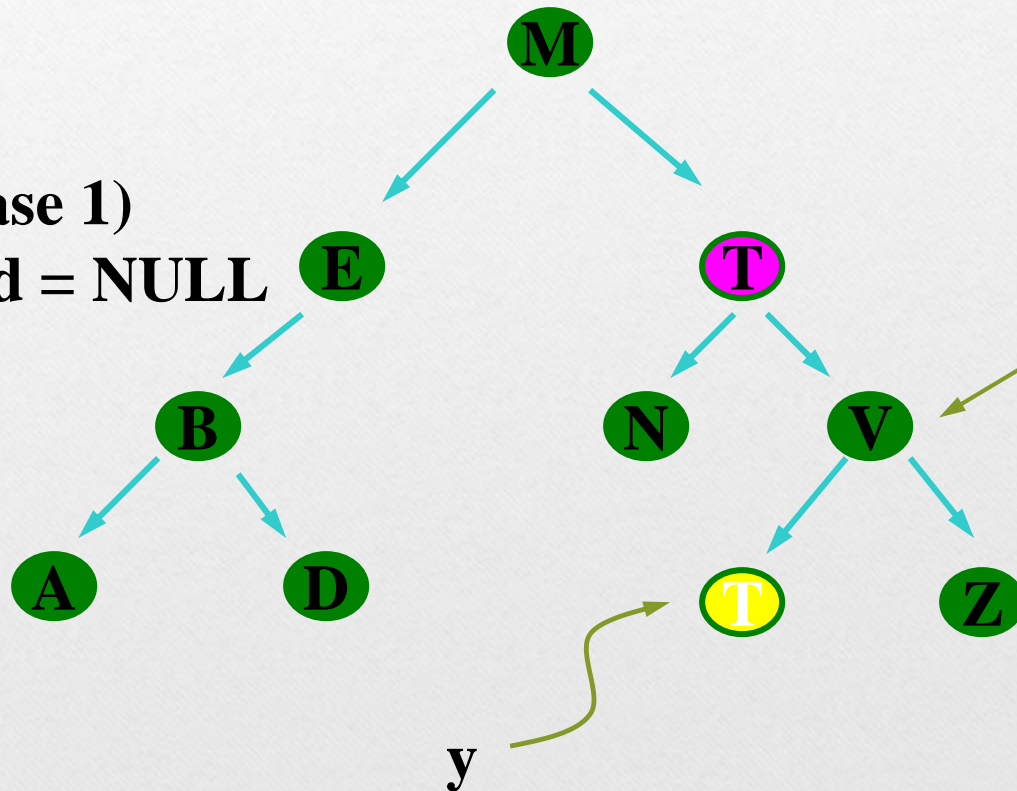
Delete Node: Third Case

parent

To delete P

Delete T (as in case 1)

parent_y->Lchild = NULL



Delete Node: Third Case

