

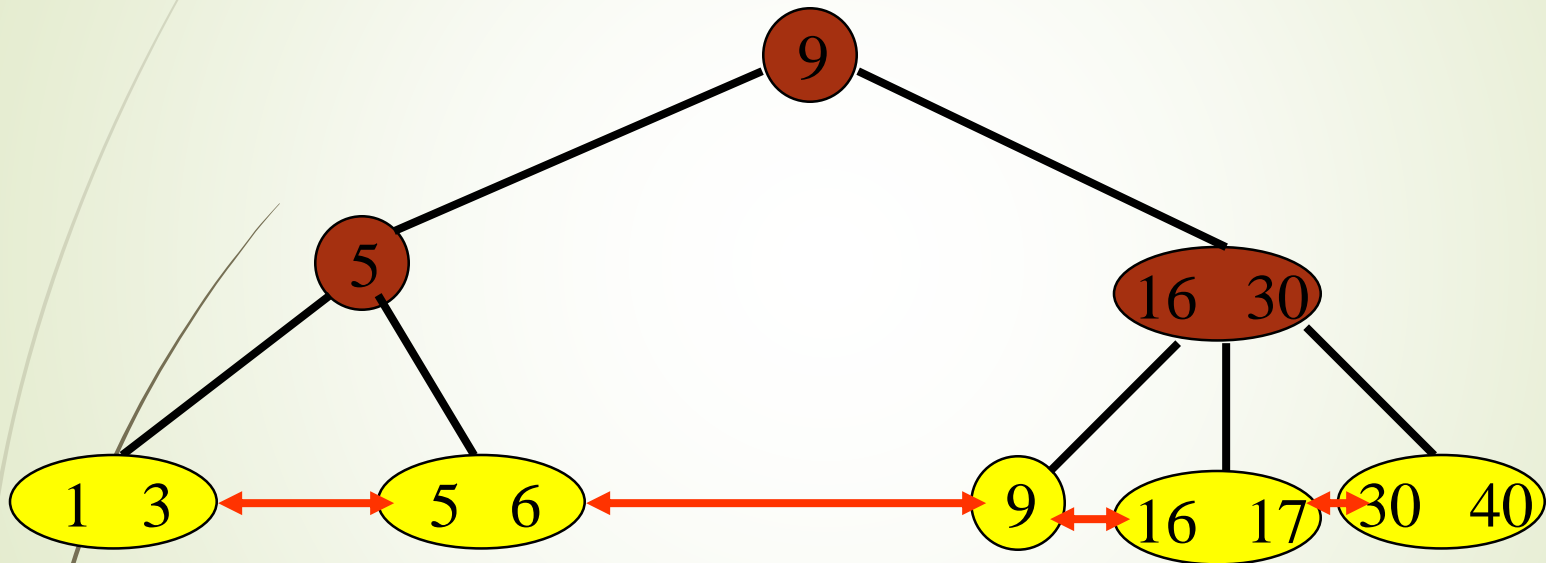
B⁺-Trees

- Same structure as B-trees.
- Dictionary pairs are in leaves only. Leaves form a doubly-linked list.
- Remaining nodes have following structure:

$j \ a_0 \ k_1 \ a_1 \ k_2 \ a_2 \ \dots \ k_j \ a_j$

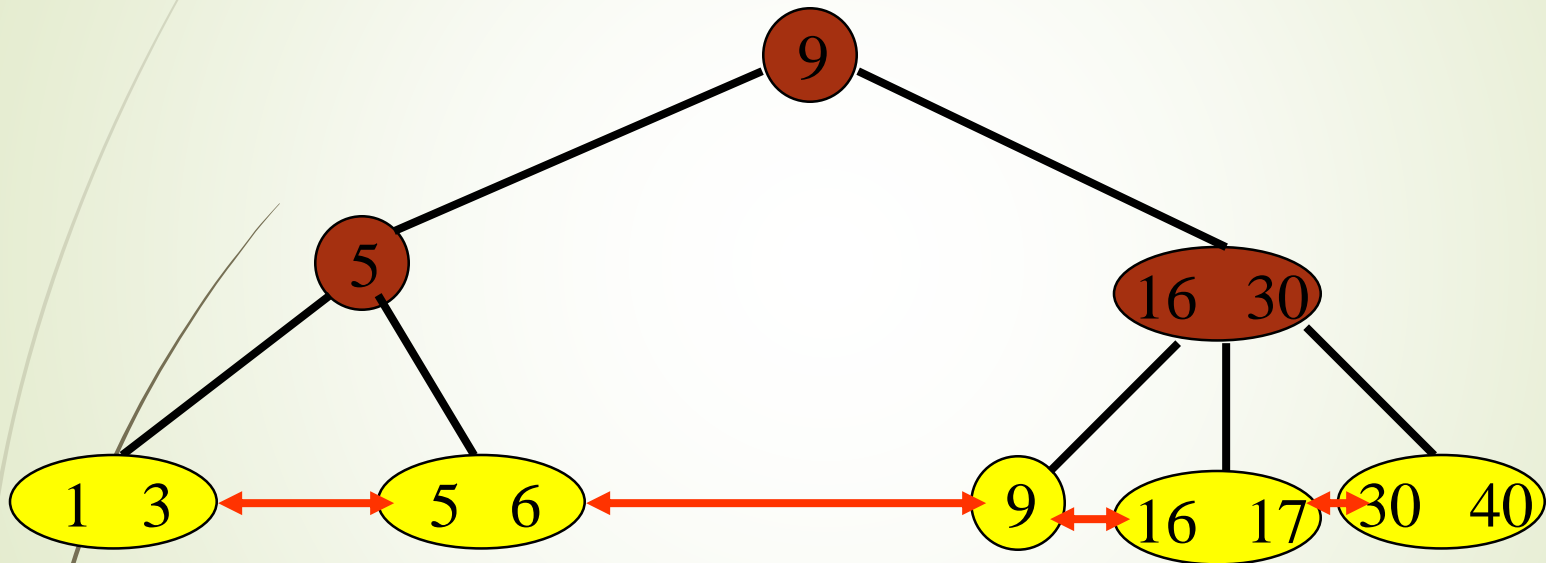
- j = number of keys in node.
- a_i is a pointer to a subtree.
- $k_i \leq$ smallest key in subtree a_i and $>$ largest in a_{i-1} .

Example B+-tree



-  → index node
-  → leaf/data node

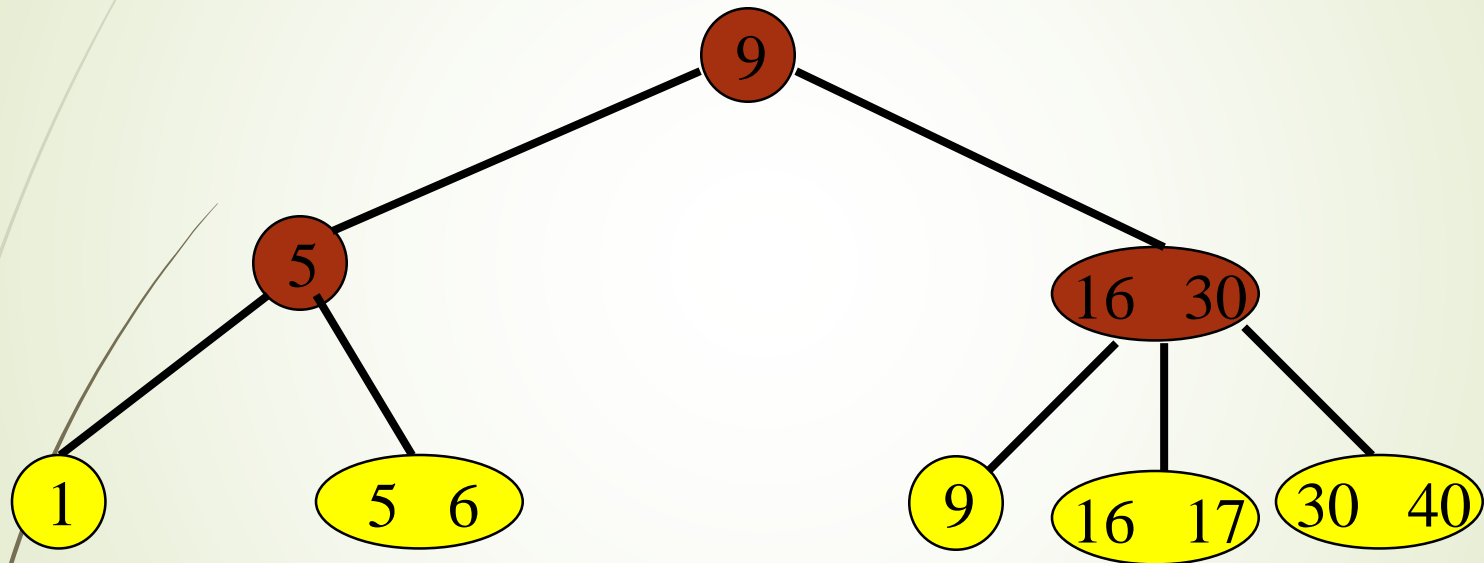
B+-tree—Search



key = 5

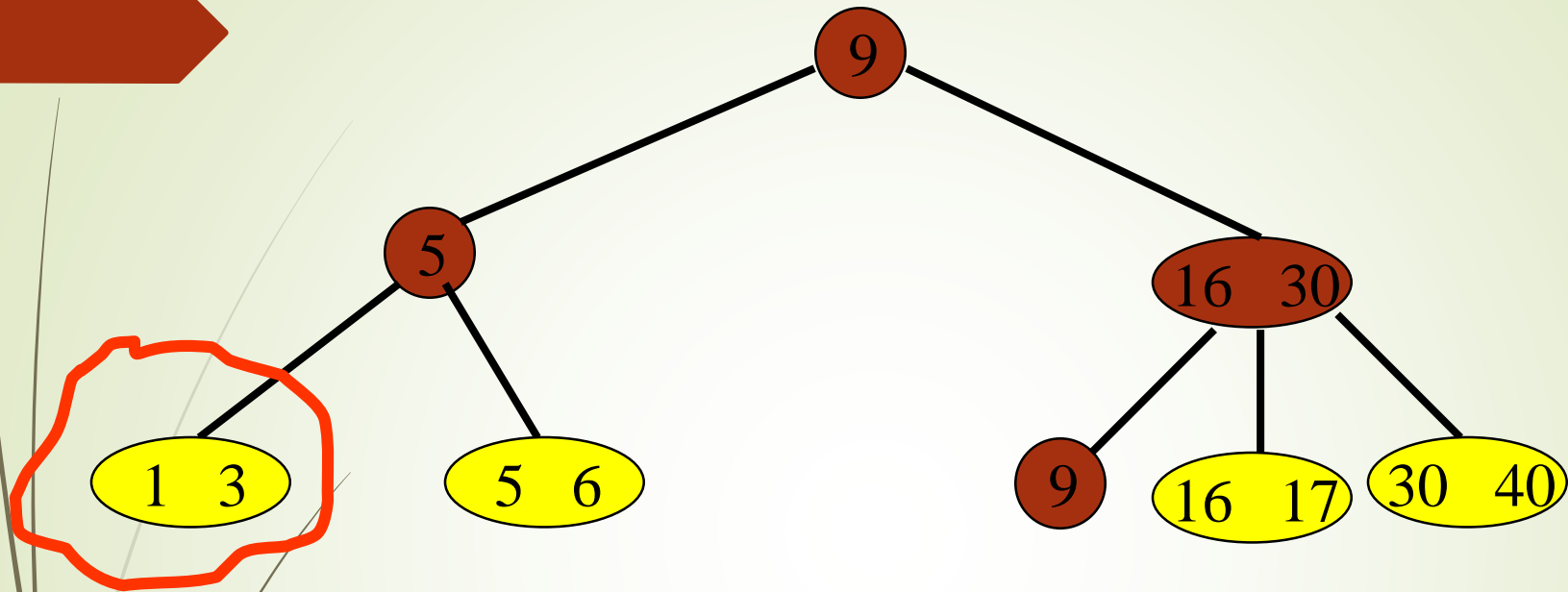
$6 \leq \text{key} \leq 20$

B+-tree—Insert



Insert 10

Insert



- Insert a pair with key = 2.
- New pair goes into a 3-node.

Insert Into A 3-node

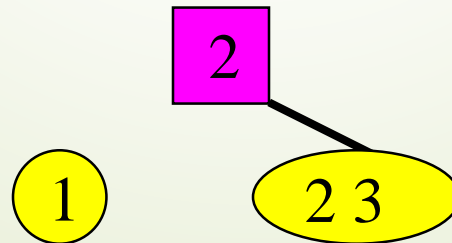
Insert new pair so that the keys are in ascending order.



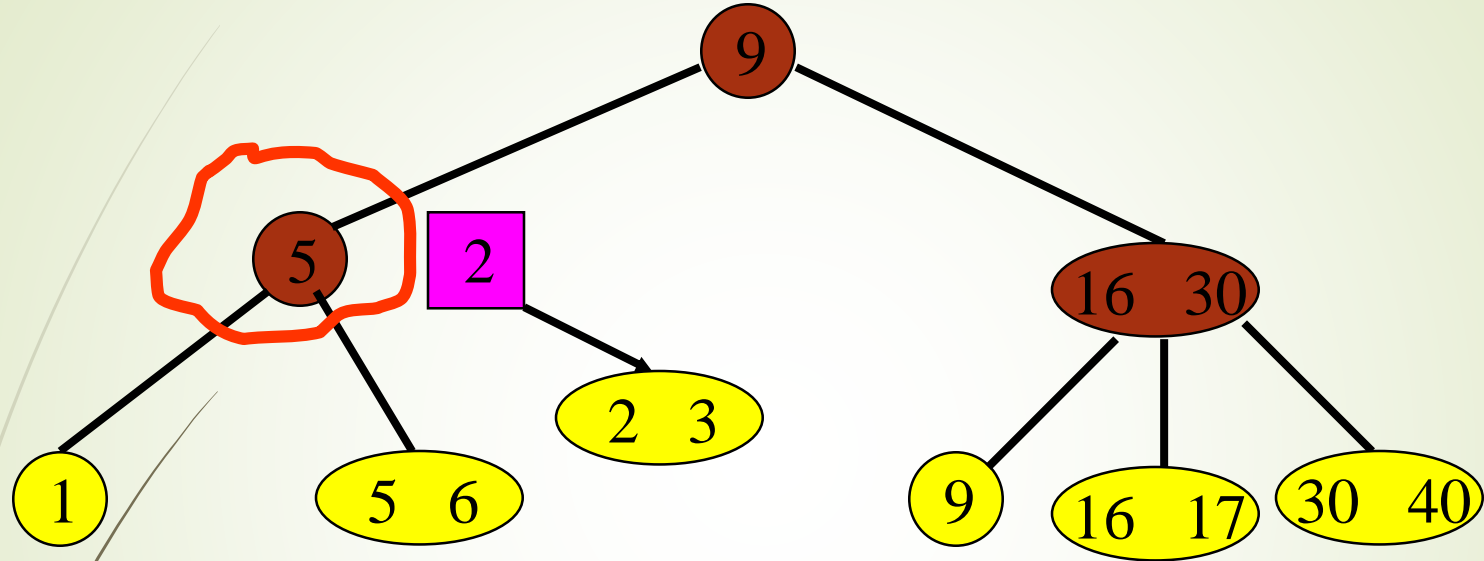
- Split into two nodes.



- Insert smallest key in new node and pointer to this new node into parent.

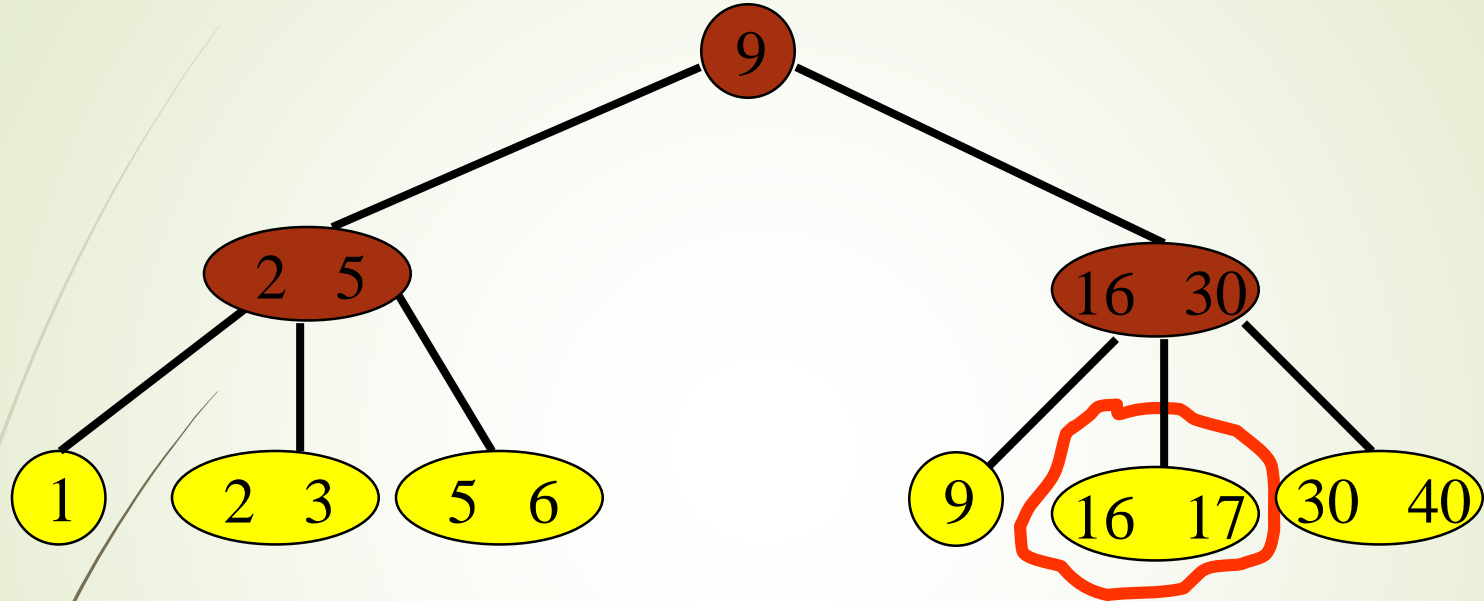


Insert



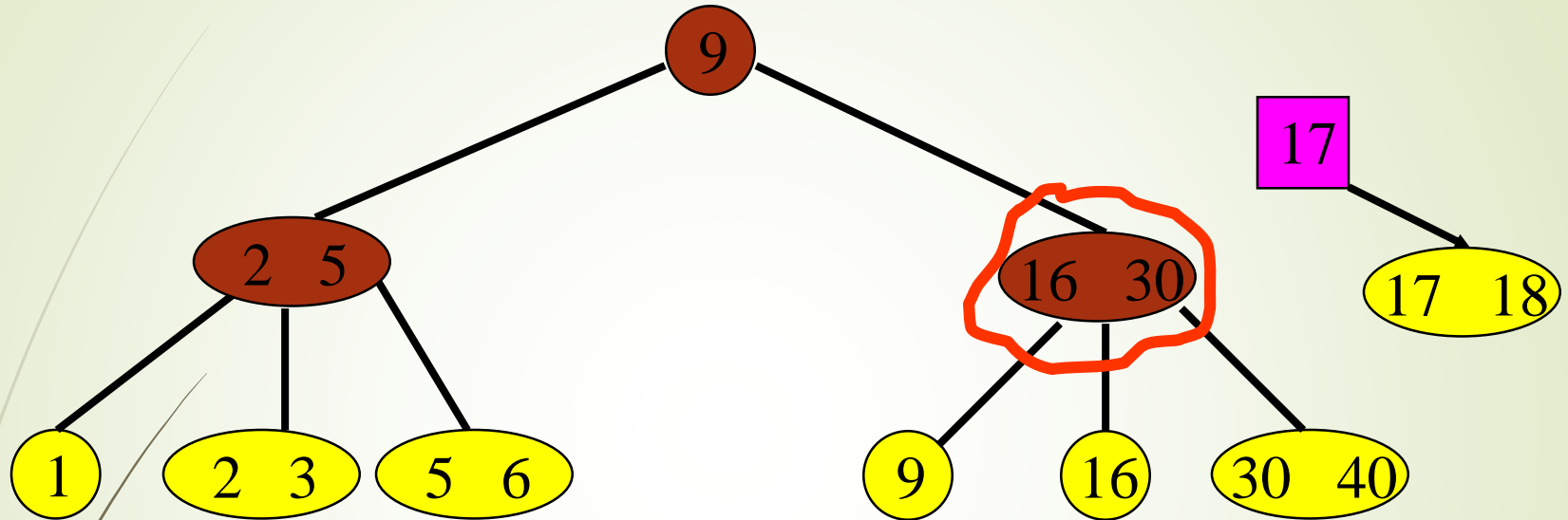
- Insert an index entry 2 plus a pointer into parent.

Insert



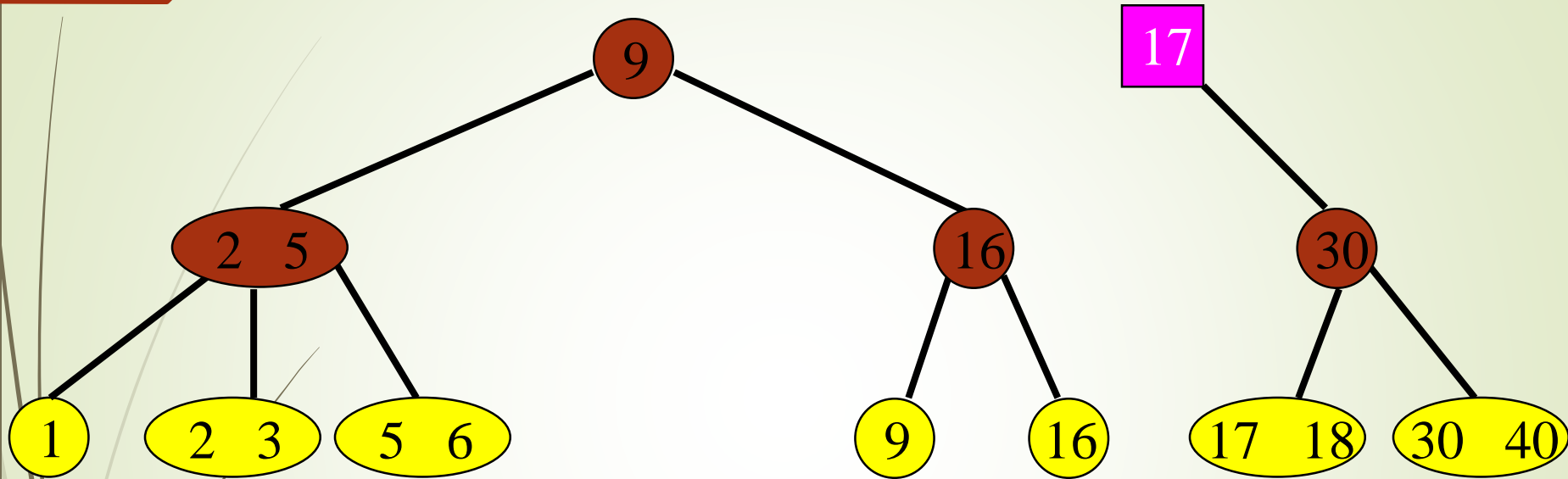
- Now, insert a pair with key = 18.

Insert



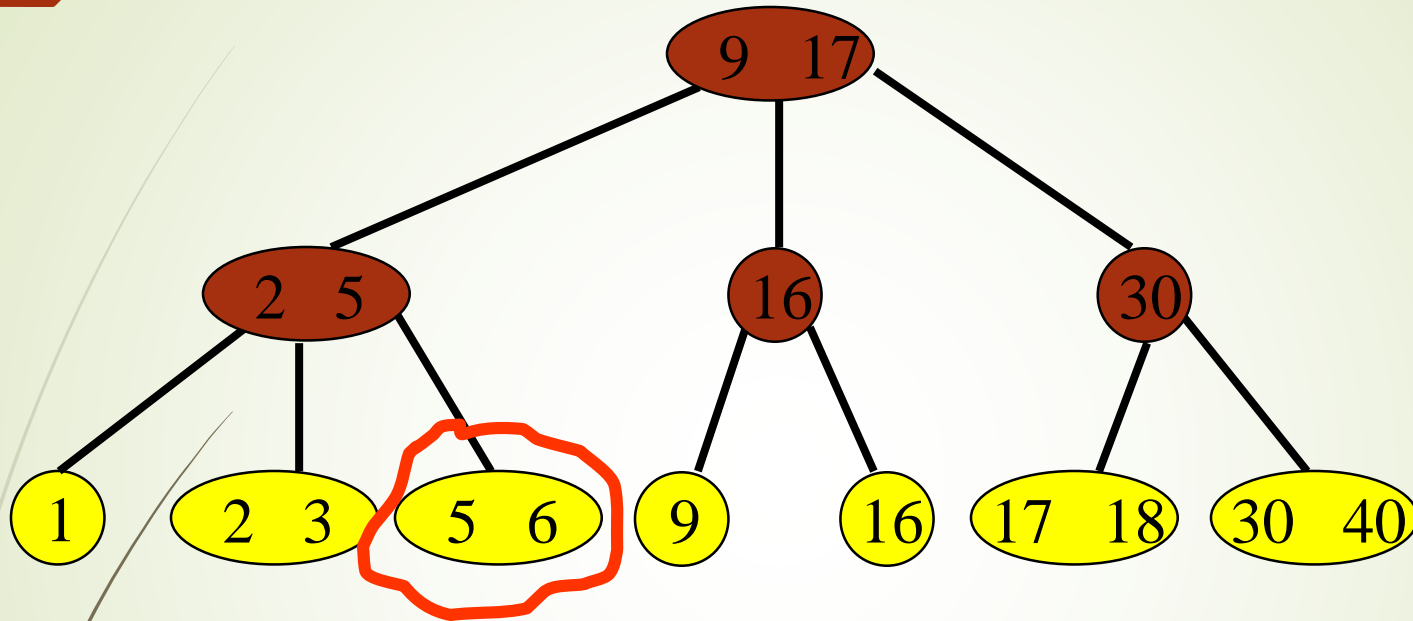
- Now, insert a pair with key = 18.
- Insert an index entry 17 plus a pointer into parent.

Insert



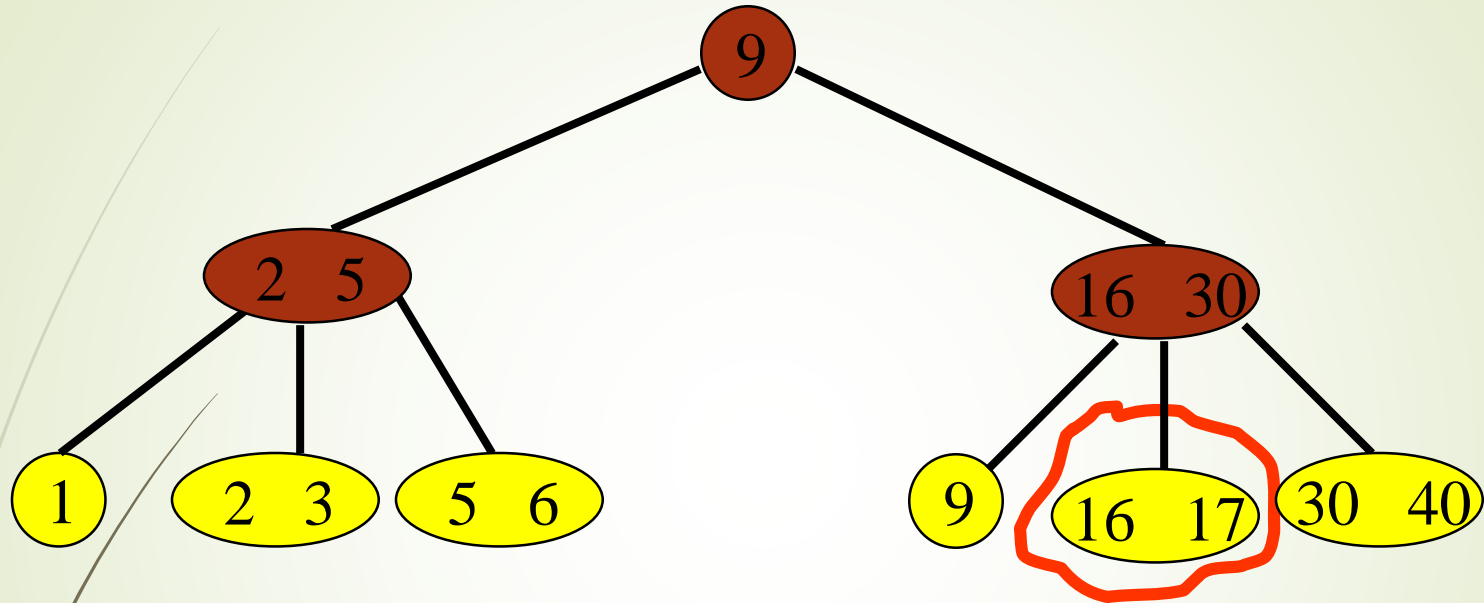
- Now, insert a pair with key = 18.
- Insert an index entry 17 plus a pointer into parent.

Insert



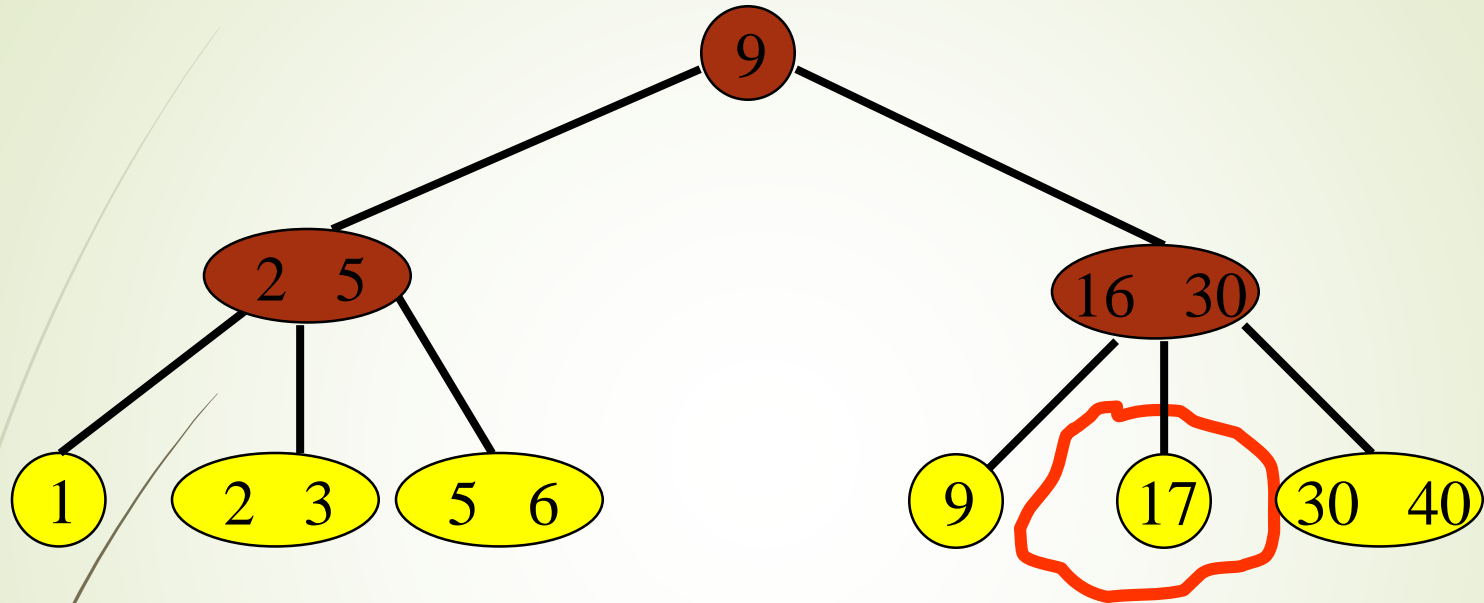
- Now, insert a pair with key = 7.

Delete



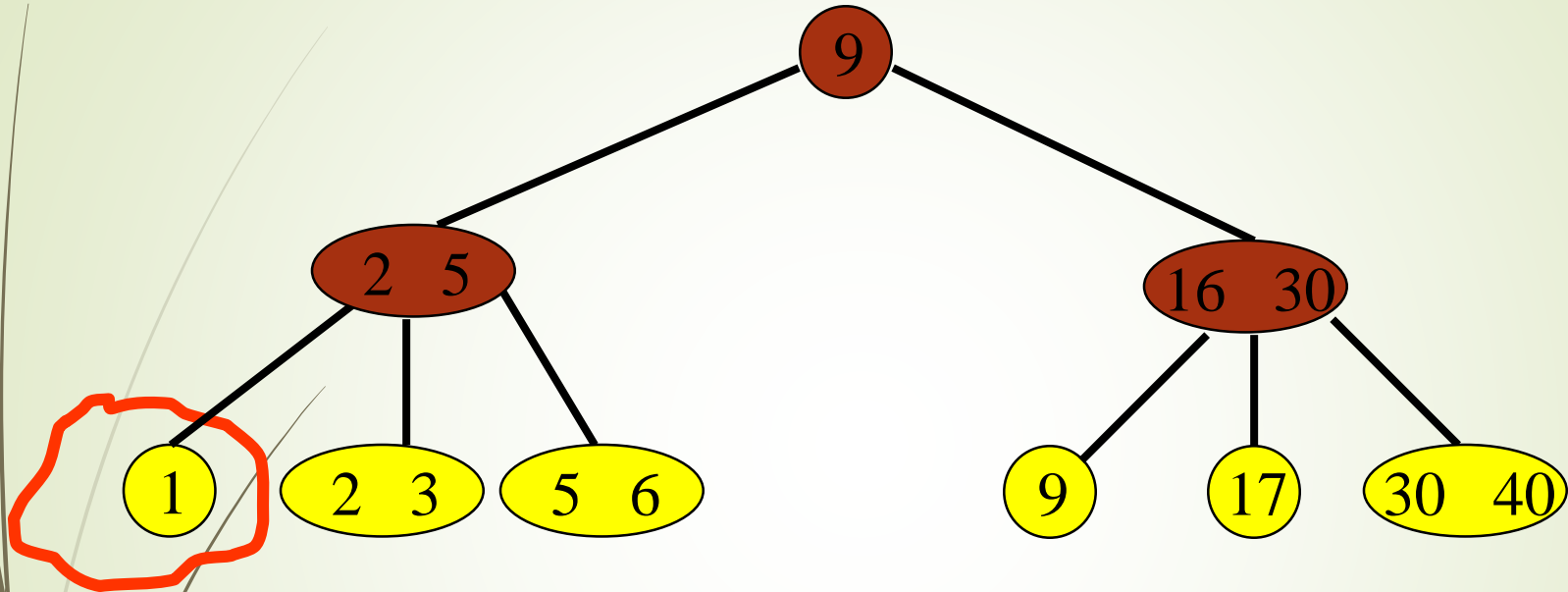
- Delete pair with key = 16.
- Note: delete pair is always in a leaf.

Delete



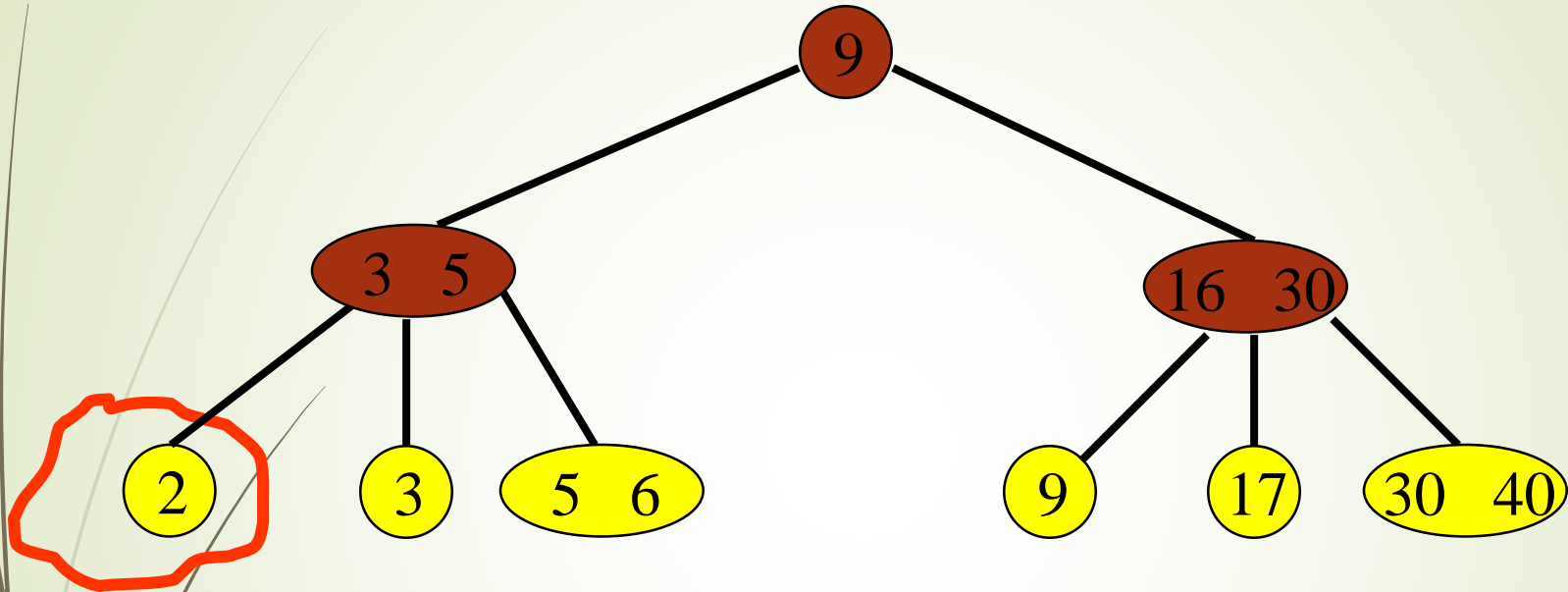
- Delete pair with key = 16.
- Note: delete pair is always in a leaf.

Delete



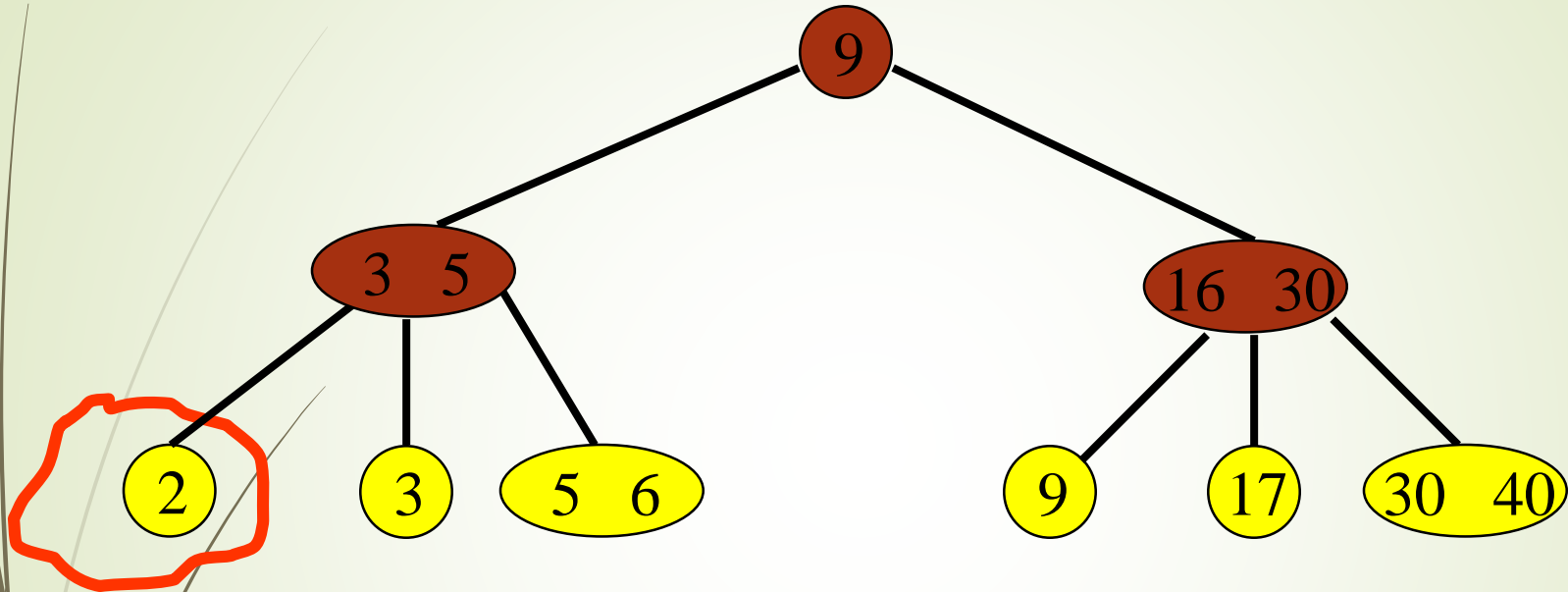
- Delete pair with key = 1.
- Get ≥ 1 from sibling and update parent key.

Delete



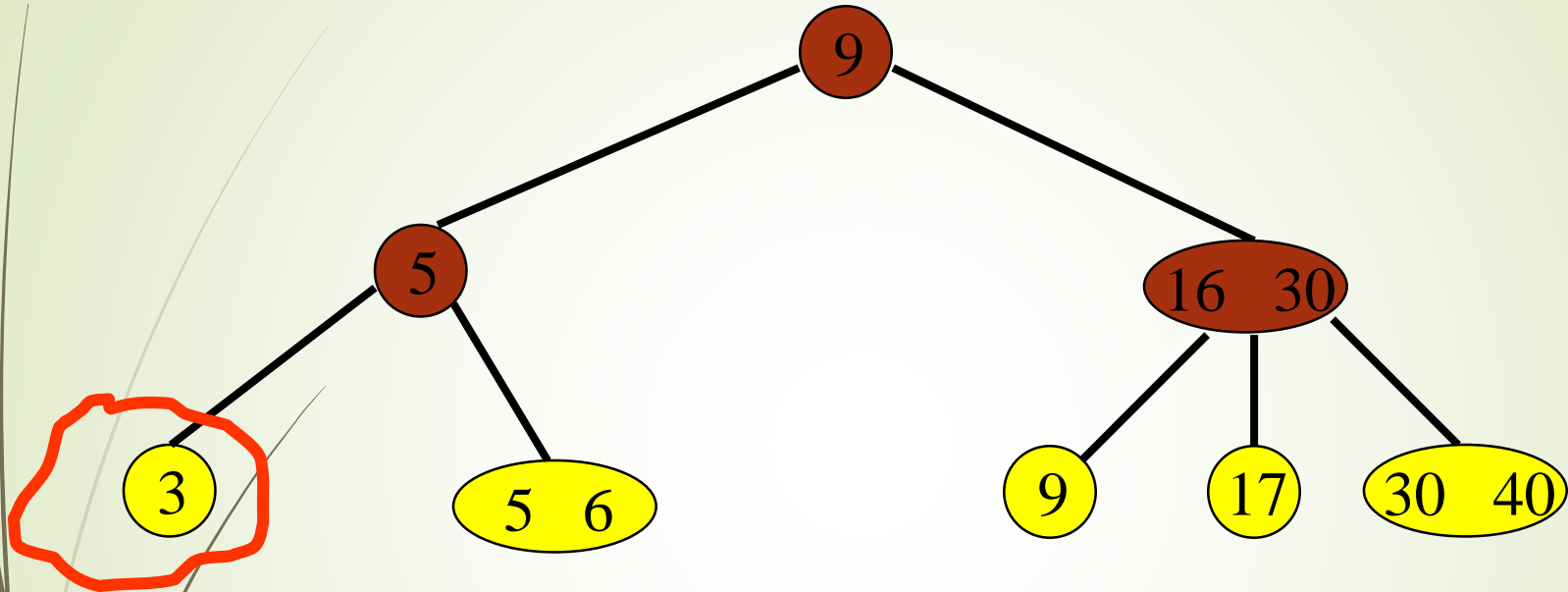
- Delete pair with key = 1.
- Get ≥ 1 from sibling and update parent key.

Delete



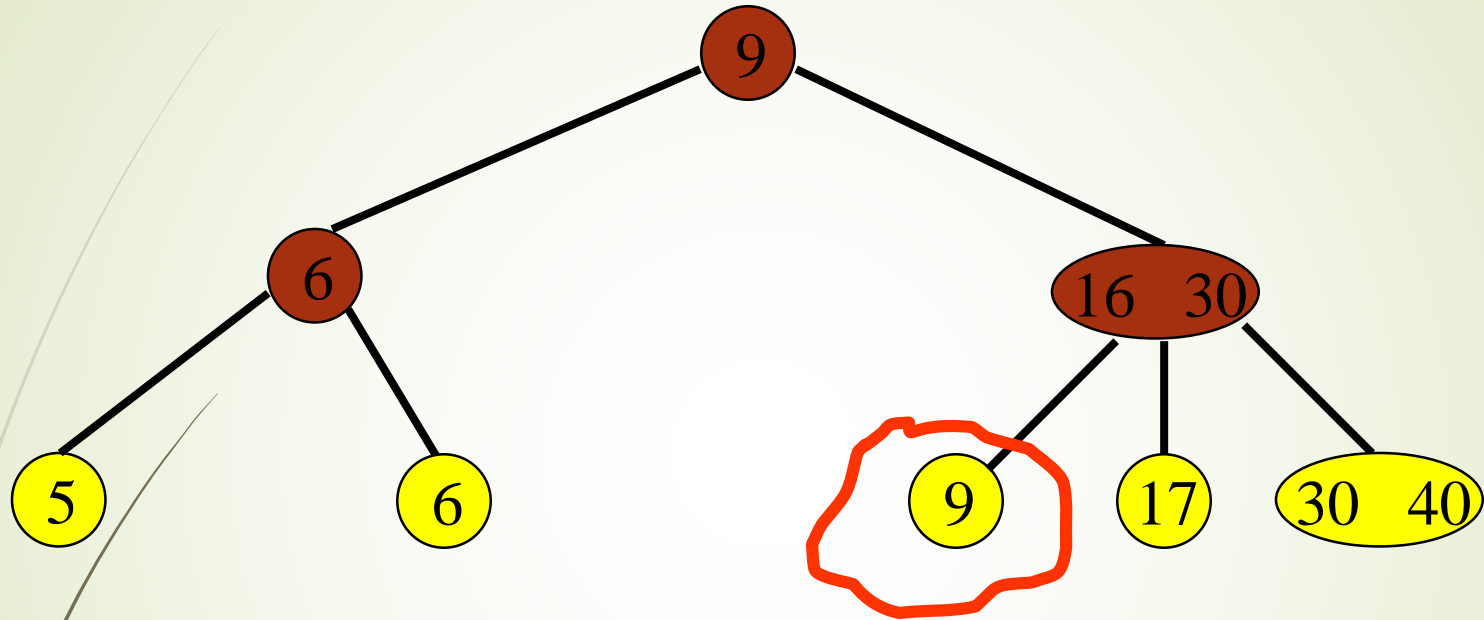
- Delete pair with key = 2.
- Merge with sibling, delete in-between key in parent.

Delete



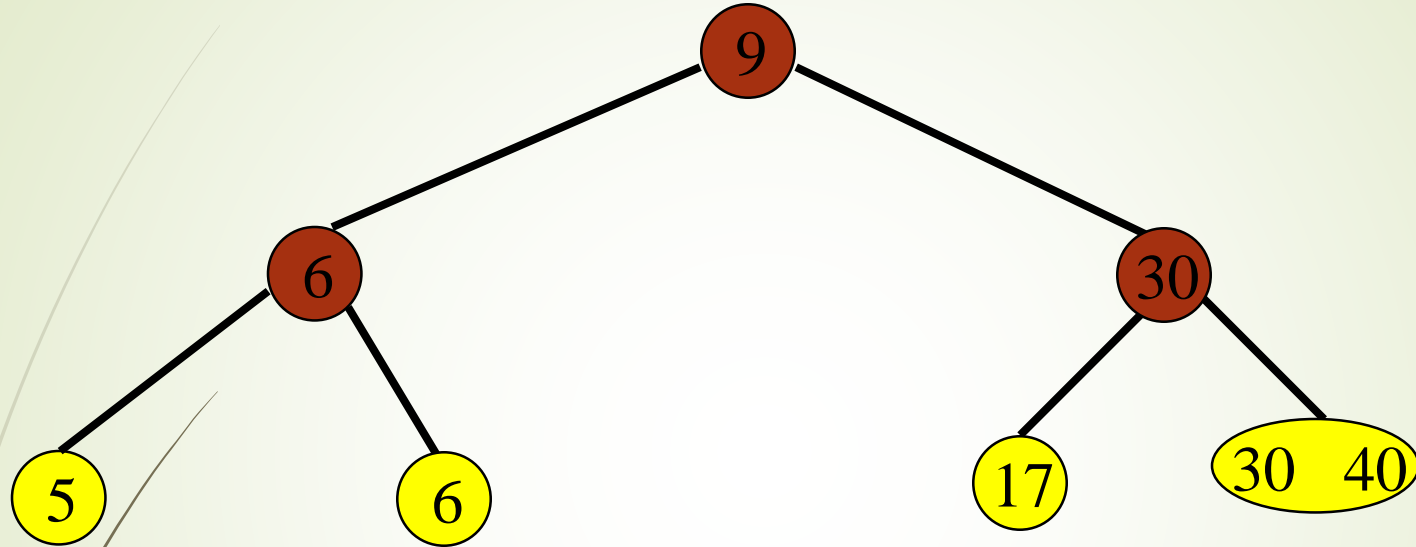
- Delete pair with key = 3.
- Get ≥ 1 from sibling and update parent key.

Delete

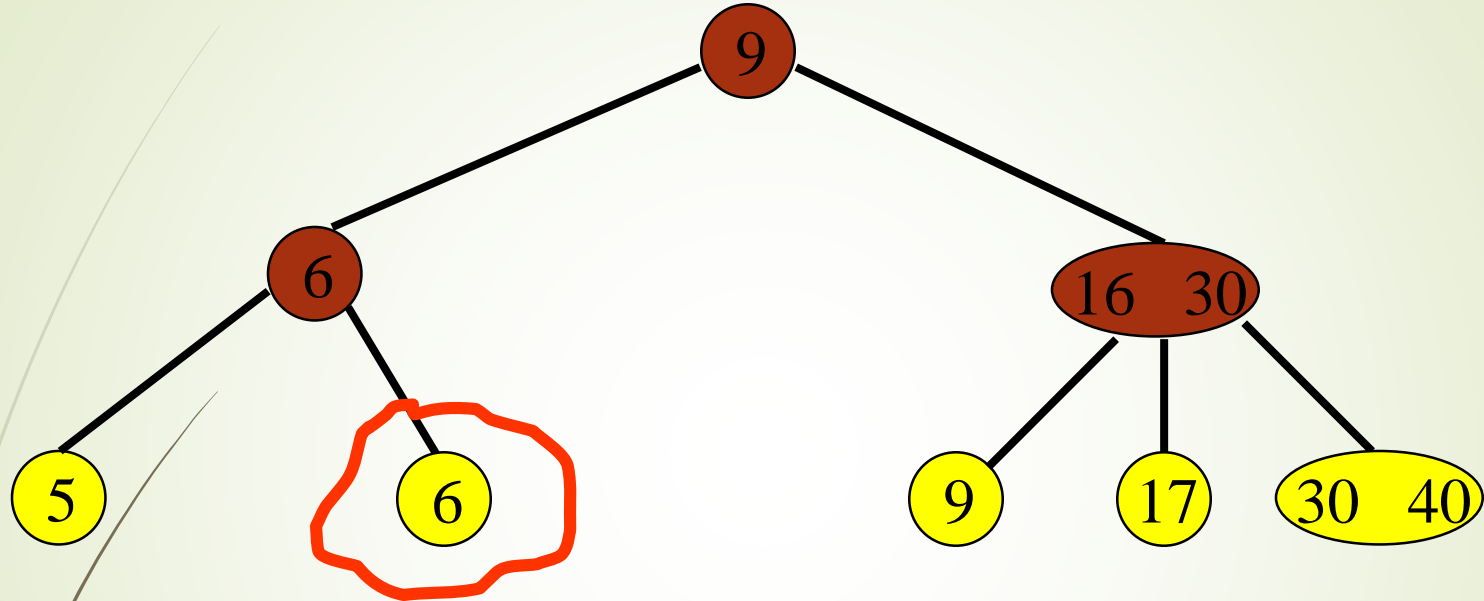


- Delete pair with key = 9.
- Merge with sibling, delete in-between key in parent.

Delete

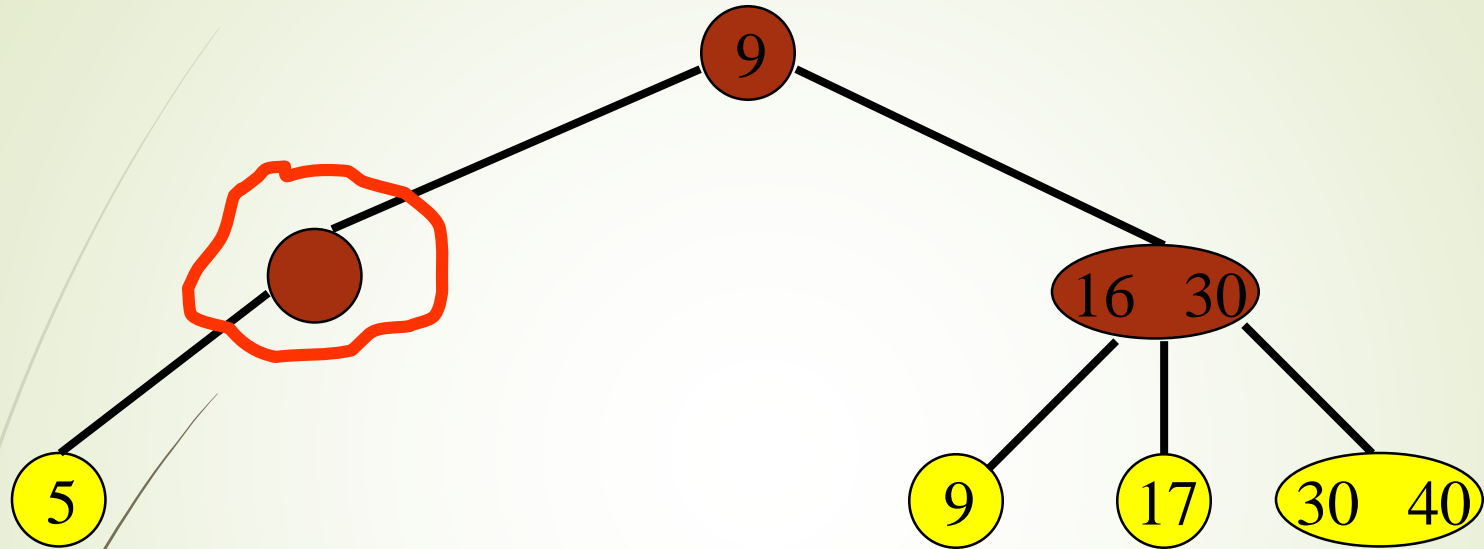


Delete



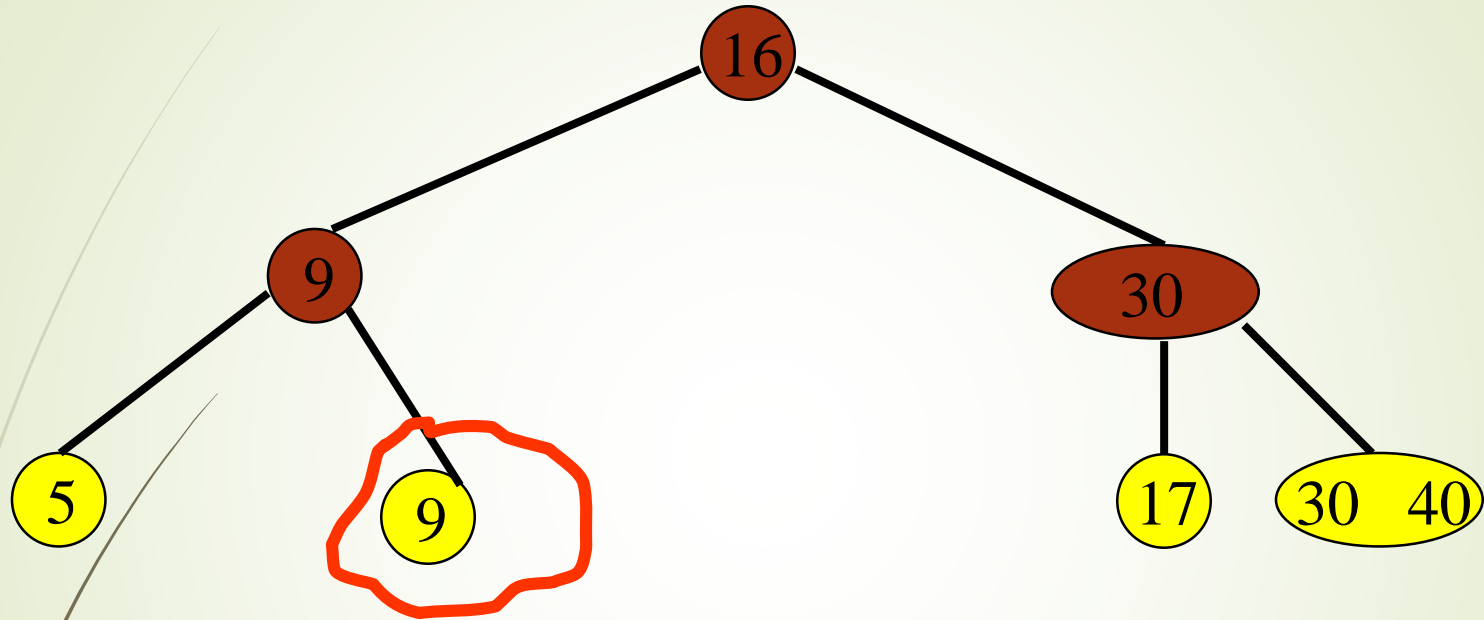
- Delete pair with key = 6.
- Merge with sibling, delete in-between key in parent.

Delete



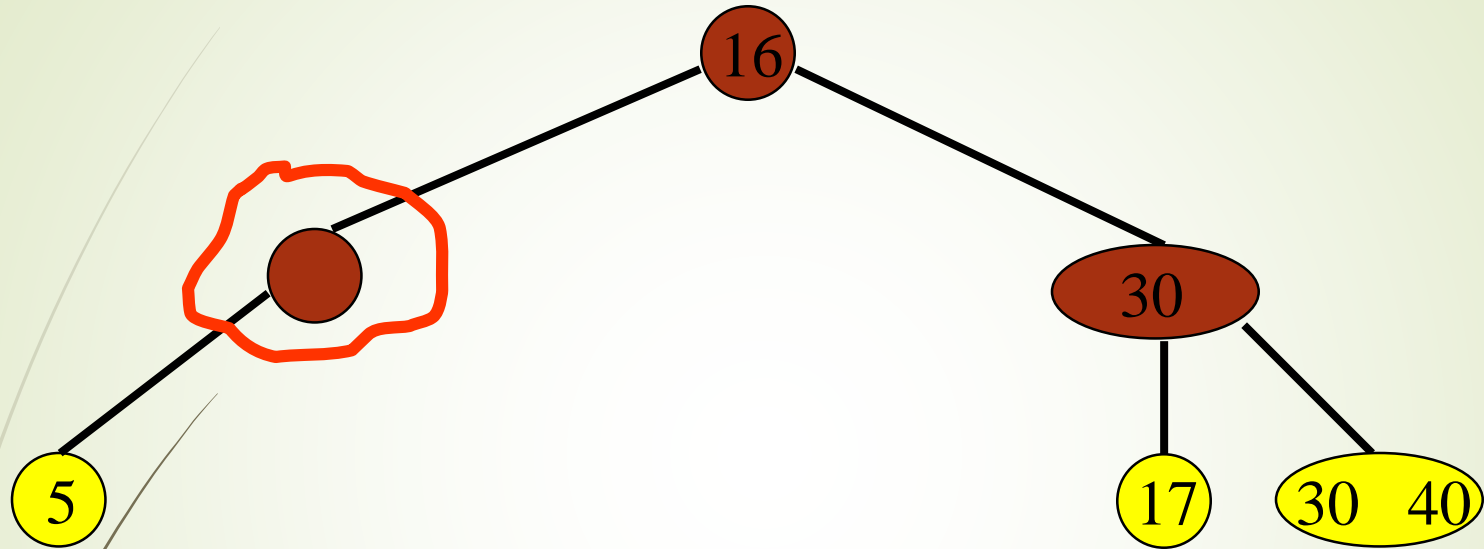
- Index node becomes deficient.
- Get ≥ 1 from sibling, move last one to parent, get parent key.

Delete



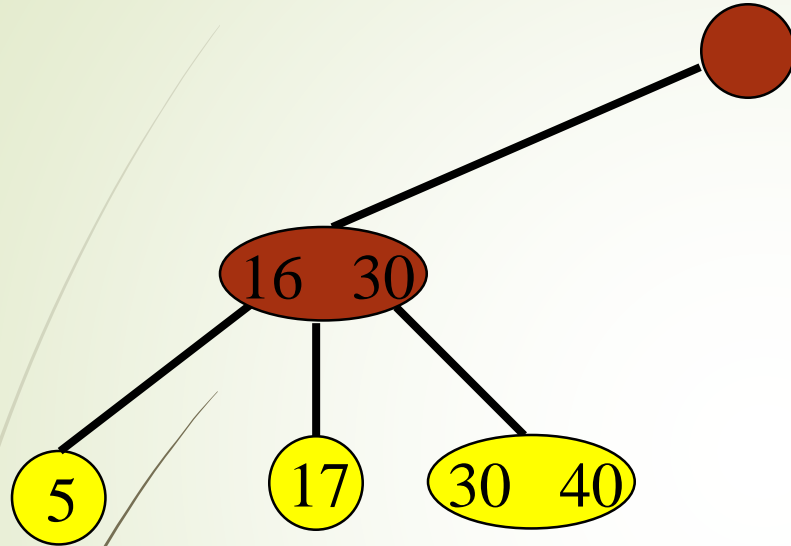
- Delete 9.
- Merge with sibling, delete in-between key in parent.

Delete



- Index node becomes deficient.
- Merge with sibling and in-between key in parent.

Delete



- Index node becomes deficient.
- It's the root; discard.

B*-Trees

- ▶ Root has between 2 and $2 * \text{floor}((2m - 2)/3) + 1$ children.
- ▶ Remaining nodes have between $\text{ceil}((2m - 1)/3)$ and m children.
- ▶ All external/failure nodes are on the same level.

Insert

- ▶ When insert node is overfull, check adjacent sibling.
- ▶ If adjacent sibling is not full, move a dictionary pair from overfull node, via parent, to nonfull adjacent sibling.
- ▶ If adjacent sibling is full, split overfull node, adjacent full node, and in-between pair from parent to get three nodes with $\text{floor}((2m - 2)/3)$, $\text{floor}((2m - 1)/3)$, $\text{floor}(2m/3)$ pairs plus two additional pairs for insertion into parent.

Delete

- ▶ When combining, must combine 3 adjacent nodes and 2 in-between pairs from parent.
 - ▶ Total # pairs involved = $2 * \text{floor}((2m-2)/3) + [\text{floor}((2m-2)/3) - 1] + 2$
 - ▶ Equals $3 * \text{floor}((2m-2)/3) + 1$.
- ▶ Combining yields 2 nodes and a pair that is to be inserted into the parent.
 - ▶ $m \bmod 3 = 0 \Rightarrow$ nodes have $m - 1$ pairs each.
 - ▶ $m \bmod 3 = 1 \Rightarrow$ one node has $m - 1$ pairs and the other has $m - 2$
 - ▶ $m \bmod 3 = 2 \Rightarrow$ nodes have $m - 2$ pairs each.